

18-660: Numerical Methods for Engineering Design and Optimization

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Overview

- Constrained Optimization
 - ▼ Linear equality constraint
 - ▼ Lagrange multiplier

Constrained Nonlinear Optimization

$$\begin{array}{ll} \min_x & f(X) \\ \text{S.T.} & \begin{cases} g_1(X) \leq 0 \\ g_2(X) \leq 0 \\ \vdots \end{cases} \end{array}$$

- Equality constraint can be written as two inequality constraints

$$\begin{array}{ll} \min_x & f(X) \\ \text{S.T.} & g(X) = 0 \end{array} \quad \Rightarrow \quad \begin{array}{ll} \min_x & f(X) \\ \text{S.T.} & \begin{cases} g(X) \leq 0 \\ -g(X) \leq 0 \end{cases} \end{array}$$

Linear Equality Constraint

$$\begin{array}{ll} \min_{X} & f(X) \\ \text{S.T.} & AX = B \end{array}$$

- Linear equality constraint can be efficiently handled by a number of optimization algorithms
 - ▼ We do not write $AX = B$ as two inequality constraints
 - ▼ It can be directly solved with high efficiency

Subspace Reduction

■ Eliminate linear equality constraint

$$\begin{array}{ccc} \begin{array}{c} \underline{AX} = \underline{B} \\ \downarrow \quad \downarrow \\ P \times N \quad P \times 1 \end{array} & \Rightarrow & \begin{array}{c} X = \underline{FZ} + D \quad (\forall Z \in R^{N-P}) \\ \swarrow \quad \searrow \\ N \times (N-P) \quad (N-P) \times 1 \end{array} \end{array}$$

- ▼ $X = FZ + D$ is the (non-unique) solution of under-determined linear equation $AX = B$
- ▼ For any Z value, AX is equal to B

A Simple Example

$$AX = B \quad \Rightarrow \quad X = FZ + D$$

$$x_1 + x_2 = 1$$

$$x_1 = z$$

$$x_2 = -z + 1$$

$$\begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ -1 \end{bmatrix} \cdot z + \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$

Subspace Reduction

$$\begin{array}{l} \min_X f(X) \\ \text{S.T. } AX = B \\ \\ X = FZ + D \end{array} \quad \Rightarrow \quad \min_Z f(FZ + D)$$

- Solve the optimal value Z by unconstrained optimization – minimizing $f(FZ+D)$
- Calculate the optimum $X = FZ+D$

Lagrange Multiplier

- Equality constraint can also be handled by Lagrange multiplier
- If X^* is a local minimum of

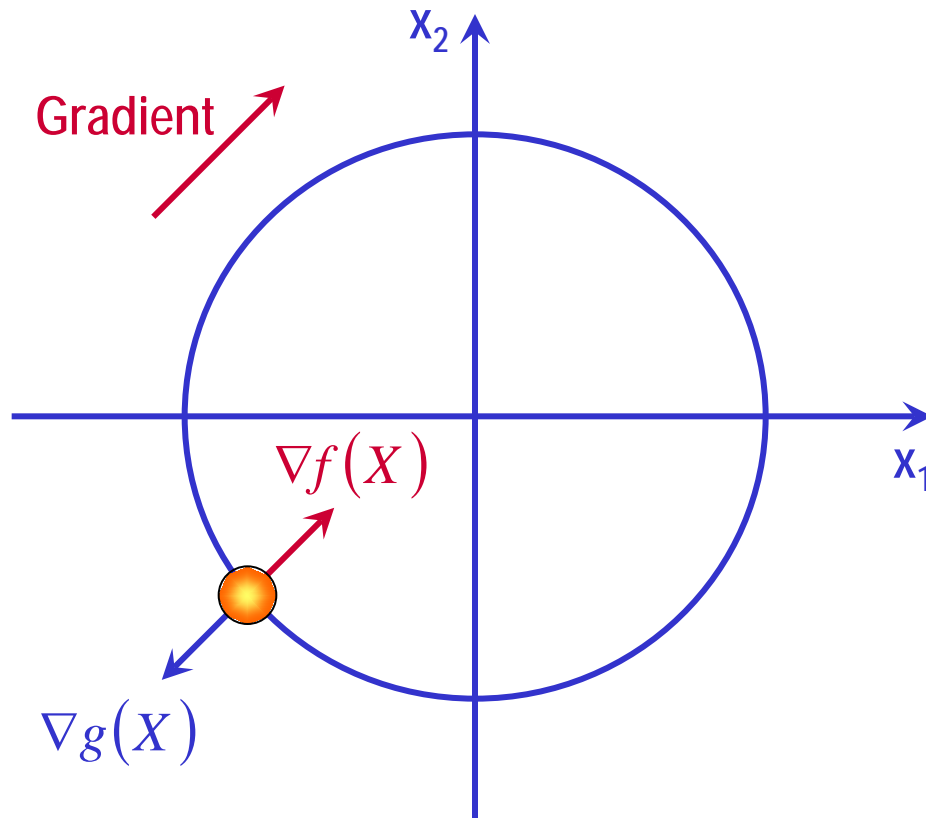
$$\begin{array}{ll} \min_X & f(X) \\ \text{S.T.} & g_i(X) = 0 \quad (i = 1, 2, \dots, P) \end{array}$$

- ▼ there exist $\lambda_1, \lambda_2, \dots, \lambda_P$, called **Lagrange multipliers**, such that

$$\nabla f(X^*) + \sum_{i=1}^P \lambda_i \cdot \nabla g_i(X^*) = 0$$

A Simple Example

$$\begin{aligned} \min_{x_1, x_2} \quad & f(X) = x_1 + x_2 \\ \text{S.T.} \quad & g(X) = x_1^2 + x_2^2 - 2 = 0 \end{aligned}$$



$$\begin{aligned} X^* &= \begin{bmatrix} -1 \\ -1 \end{bmatrix} \\ \nabla f(X^*) &= \begin{bmatrix} 1 \\ 1 \end{bmatrix} \\ \nabla g(X^*) &= \begin{bmatrix} -2 \\ -2 \end{bmatrix} \end{aligned}$$

Lagrange Multiplier

$$\begin{array}{ll} \min_X & f(X) \\ \text{S.T.} & g_i(X) = 0 \quad (i = 1, 2, \dots, P) \end{array} \quad \Rightarrow \quad \nabla f(X^*) + \sum_{i=1}^P \lambda_i \cdot \nabla g_i(X^*) = 0$$

■ Optimality condition for linear constraints

$$\begin{array}{ll} \min_X & f(X) \\ \text{S.T.} & AX = B \end{array}$$

$$g_i(X) = A(i,:) \cdot X - B(i) = 0$$

$$\left[\begin{array}{c} \text{---} \\ A \\ \text{---} \end{array} \right] \cdot X - \left[\begin{array}{c} \text{---} \\ B \\ \text{---} \end{array} \right] = 0 \quad \rightarrow \text{i-th row}$$

$$\nabla g_i(X) = A(i,:)^T$$

Lagrange Multiplier

$$\begin{array}{ll} \min_X & f(X) \\ \text{S.T.} & g_i(X) = 0 \quad (i = 1, 2, \dots, P) \end{array} \quad \Rightarrow \quad \nabla f(X^*) + \sum_{i=1}^P \lambda_i \cdot \nabla g_i(X^*) = 0$$

■ Optimality condition for linear constraints

$$\nabla g_i(X) = A(i, :)^T$$

$$\sum_{i=1}^P \lambda_i \cdot \nabla g_i(X) = \sum_{i=1}^P \lambda_i \cdot A(i, :)^T = A^T V$$

$$\nabla f(X^*) + A^T V = 0$$

$$\underbrace{\begin{bmatrix} A(1, :)^T & A(2, :)^T & \dots \end{bmatrix}}_{A^T} \cdot \underbrace{\begin{bmatrix} \lambda_1 \\ \lambda_2 \\ \vdots \end{bmatrix}}_V$$

Linear Equality Constrained Quadratic Programming

$$\begin{array}{ll} \min_X & f(X) \\ \text{S.T.} & AX = B \end{array} \quad \Rightarrow \quad \nabla f(X^*) + A^T V = 0$$

■ We first consider quadratic cost function

- ▼ Any smooth nonlinear cost function can be locally approximated as a quadratic function (2nd-order Taylor expansion)

$$\begin{array}{ll} \min_X & f(X) = \frac{1}{2} X^T Q X + R^T X + C \\ \text{S.T.} & AX = B \end{array}$$

$$\nabla f(X) = QX + R$$

$$QX^* + R + A^T V = 0$$

Linear Equality Constrained Quadratic Programming

$$\begin{aligned} \min_x \quad & \frac{1}{2} X^T Q X + R^T X + C \\ \text{S.T.} \quad & A X = B \end{aligned}$$

- Optimality condition for quadratic programming

$$Q X^* + R + A^T V = 0$$

$$A X^* = B$$

$$\begin{bmatrix} Q & A^T \\ A & 0 \end{bmatrix} \cdot \begin{bmatrix} X^* \\ V \end{bmatrix} = \begin{bmatrix} -R \\ B \end{bmatrix}$$

Solve linear equation to determine X^* (optimal solution) and V (Lagrange multipliers)

Linear Equality Constrained Nonlinear Programming

$$\begin{array}{ll} \min_X & \frac{1}{2} X^T Q X + R^T X + C \\ \text{S.T.} & A X = B \end{array} \quad \Rightarrow \quad \begin{bmatrix} Q & A^T \\ A & 0 \end{bmatrix} \cdot \begin{bmatrix} X^* \\ V \end{bmatrix} = \begin{bmatrix} -R \\ B \end{bmatrix}$$

- Minimize nonlinear function $f(X)$ given linear constraint $A X = B$

$$\begin{array}{ll} \min_X & f(X) \\ \text{S.T.} & A X = B \end{array}$$

$$f[X^{(k+1)}] \approx \frac{1}{2} \cdot \Delta X^T \cdot \nabla^2 f[X^{(k)}] \cdot \Delta X + \nabla f[X^{(k)}]^T \cdot \Delta X + f[X^{(k)}]$$

$$A \cdot \Delta X = A \cdot [X^{(k+1)} - X^{(k)}] = A \cdot X^{(k+1)} - A X^{(k)} = B - A X^{(k)}$$

Linear Equality Constrained Nonlinear Programming

$$\begin{array}{ll} \min_X & \frac{1}{2} X^T Q X + R^T X + C \\ \text{S.T.} & A X = B \end{array} \quad \Rightarrow \quad \begin{bmatrix} Q & A^T \\ A & 0 \end{bmatrix} \cdot \begin{bmatrix} X^* \\ V \end{bmatrix} = \begin{bmatrix} -R \\ B \end{bmatrix}$$

$$f[X^{(k+1)}] \approx \frac{1}{2} \cdot \Delta X^T \cdot \nabla^2 f[X^{(k)}] \cdot \Delta X + \nabla f[X^{(k)}]^T \cdot \Delta X + f[X^{(k)}]$$

$$A \cdot \Delta X = B - A X^{(k)}$$

$$\begin{bmatrix} \nabla^2 f[X^{(k)}] & A^T \\ A & 0 \end{bmatrix} \cdot \begin{bmatrix} \Delta X \\ V \end{bmatrix} = \begin{bmatrix} -\nabla f[X^{(k)}] \\ B - A X^{(k)} \end{bmatrix}$$

Solve linear equation to determine ΔX

$$X^{(k+1)} = X^{(k)} + \Delta X$$

Update $X^{(k+1)}$

Linear Equality Constrained Nonlinear Programming

$$\begin{array}{ll} \min_X & f(X) \\ \text{S.T.} & AX = B \end{array} \quad \Rightarrow \quad \begin{bmatrix} \nabla^2 f[X^{(k)}] & A^T \\ A & 0 \end{bmatrix} \cdot \begin{bmatrix} \Delta X \\ V \end{bmatrix} = \begin{bmatrix} -\nabla f[X^{(k)}] \\ B - AX^{(k)} \end{bmatrix}$$

■ If $X^{(k)}$ is a feasible solution

- ▼ We can start from an initial solution $X^{(0)}$ that is feasible
- ▼ Even if $X^{(0)}$ is not feasible, $X^{(1)}$ is feasible after one iteration

$$AX^{(k)} = B$$

$$B - AX^{(k)} = 0$$

$$\begin{bmatrix} \nabla^2 f[X^{(k)}] & A^T \\ A & 0 \end{bmatrix} \cdot \begin{bmatrix} \Delta X \\ V \end{bmatrix} = \begin{bmatrix} -\nabla f[X^{(k)}] \\ 0 \end{bmatrix}$$

A Simple Example

$$\begin{array}{ll} \min_{x_1, x_2} & x_1^4 + x_2^4 \\ \text{S.T.} & x_1 + x_2 = 1 \end{array}$$

$$X^{(0)} = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

(Feasible solution)

$$\begin{array}{ll} \min_X & f(X) \\ \text{S.T.} & AX = B \end{array}$$

$$f(X) = x_1^4 + x_2^4$$

$$A = \begin{bmatrix} 1 & 1 \end{bmatrix}$$

$$B = 1$$

$$\nabla^2 f(X) = \begin{bmatrix} 12x_1^2 & 0 \\ 0 & 12x_2^2 \end{bmatrix}$$

$$\nabla f(X) = \begin{bmatrix} 4x_1^3 \\ 4x_2^3 \end{bmatrix}$$

A Simple Example

$$\begin{array}{ll} \min_{x_1, x_2} & x_1^4 + x_2^4 \\ \text{S.T.} & x_1 + x_2 = 1 \end{array}$$

$$X^{(0)} = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$A = [1 \quad 1]$$

$$\nabla^2 f(X) = \begin{bmatrix} 12x_1^2 & 0 \\ 0 & 12x_2^2 \end{bmatrix}$$

$$\nabla f(X) = \begin{bmatrix} 4x_1^3 \\ 4x_2^3 \end{bmatrix}$$

$$\begin{bmatrix} \nabla^2 f[X^{(k)}] & A^T \\ A & 0 \end{bmatrix} \cdot \begin{bmatrix} \Delta X \\ V \end{bmatrix} = \begin{bmatrix} -\nabla f[X^{(k)}] \\ 0 \end{bmatrix}$$

$$\nabla^2 f[X^{(0)}] = \begin{bmatrix} 12 & 0 \\ 0 & 0 \end{bmatrix} \quad \nabla f[X^{(0)}] = \begin{bmatrix} 4 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} 12 & 0 & 1 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} \Delta X \\ V \end{bmatrix} = \begin{bmatrix} -4 \\ 0 \\ 0 \end{bmatrix}$$

$$\Delta X = \begin{bmatrix} -0.33 \\ 0.33 \end{bmatrix} \quad X^{(1)} = \begin{bmatrix} 0.67 \\ 0.33 \end{bmatrix}$$

A Simple Example

$$\begin{array}{ll} \min_{x_1, x_2} & x_1^4 + x_2^4 \\ \text{S.T.} & x_1 + x_2 = 1 \end{array}$$

$$X^{(1)} = \begin{bmatrix} 0.67 \\ 0.33 \end{bmatrix}$$

$$A = \begin{bmatrix} 1 & 1 \end{bmatrix}$$

$$\nabla^2 f(X) = \begin{bmatrix} 12x_1^2 & 0 \\ 0 & 12x_2^2 \end{bmatrix}$$

$$\nabla f(X) = \begin{bmatrix} 4x_1^3 \\ 4x_2^3 \end{bmatrix}$$

$$\begin{bmatrix} \nabla^2 f[X^{(k)}] & A^T \\ A & 0 \end{bmatrix} \cdot \begin{bmatrix} \Delta X \\ V \end{bmatrix} = \begin{bmatrix} -\nabla f[X^{(k)}] \\ 0 \end{bmatrix}$$

$$\nabla^2 f[X^{(1)}] = \begin{bmatrix} 5.39 & 0 \\ 0 & 1.31 \end{bmatrix} \quad \nabla f[X^{(1)}] = \begin{bmatrix} 1.20 \\ 0.14 \end{bmatrix}$$

$$\begin{bmatrix} 5.39 & 0 & 1 \\ 0 & 1.31 & 1 \\ 1 & 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} \Delta X \\ V \end{bmatrix} = \begin{bmatrix} -1.20 \\ -0.14 \\ 0 \end{bmatrix}$$

$$\Delta X = \begin{bmatrix} -0.16 \\ 0.16 \end{bmatrix} \quad X^{(2)} = \begin{bmatrix} 0.51 \\ 0.49 \end{bmatrix}$$

Linear Equality Constrained Nonlinear Programming

- Linear equality constraints can be efficiently handled by subspace reduction or Lagrange multiplier
- Nonlinear equality constraints and inequality constraints must be handled by a different algorithm
 - ▼ Interior point method (also referred to as barrier method)
 - ▼ More details in future lectures

Summary

- Constrained optimization
 - ▼ Linear equality constraint
 - ▼ Lagrange multiplier