# Interval Arithmetic FFT for Large Integer Multiplication 

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## Motivation

Large integer multiplication has applications in:

- Modular arithmetic routines in cryptographic systems
- Number theoretic tasks in kernels
- Computer algebra systems


Computation complexity of naive multiplication: $\mathrm{O}\left(\mathrm{N}^{2}\right)$ Too slow and infeasible for these applications

## Our Approach

## Fast Fourier Transform (FFT) + Interval Arithmetic

- We perform our multiplication pipeline with FFT $\rightarrow$ computation complexity $\mathrm{O}(\mathrm{N} \log \mathrm{N})$


$$
\begin{aligned}
& y[n]= \sum_{\cdots-n}^{N-1} x_{1}[m] x_{2}\left[(n-m)_{N}\right] \stackrel{\mathrm{DFT} / \mathrm{IDFT}}{\Longleftrightarrow} Y[k]=X_{1}[k] \circ X_{2}[k] \\
& 12 \times 21=[1,2] *[2,1]=[2,5,2] \\
& 12 \times 21=\mathcal{F}^{-1}(\mathcal{F}[1,2] \odot \mathcal{F}[2,1])=[2,5,2]
\end{aligned}
$$

- Pitfalls $\triangle$ : Complex FFTs lead to round-off errors associated with floating point arithmetic

|  |  |
| :---: | :---: |
|  |  |
|  |  |
|  |  |

$0 \quad \beta^{c_{\min }} \quad \beta^{c_{\min }+1} \quad \beta^{c_{\min }+2} \quad \beta^{c_{\min }+3}$

- Interval Arithmetic to handle floating point round-off error $\nabla$

$$
\left[x_{1}, x_{2}\right]+\left[y_{1}, y_{2}\right]=\left[x_{1}+y_{1}, x_{2}+y_{2}\right]
$$

$\left[x_{1}, x_{2}\right] \cdot\left[y_{1}, y_{2}\right]=\left[\min \left\{x_{1} y_{1}, x_{1} y_{2}, x_{2} y_{1}, x_{2} y_{2}\right\}, \max \left\{x_{1} y_{1}, x_{1} y_{2}, x_{2} y_{1}, x_{2} y_{2}\right\}\right]$

| Algorithm Interval-arithmetic FFT-based Integer Multiplication |  |
| :---: | :---: |
| Input: vector $x$ and $y$ of size N |  |
| Output: vector $z$ of size 2 N |  |
| $i 1 \leftarrow \operatorname{ZeroPad}(x)$ |  |
| $i 2 \leftarrow \operatorname{ZeroPad}(y)$ | $\triangleright$ zero padding to size of 2 N |
| $f 1 \leftarrow$ IntervalFFT $(i 1)$ |  |
| $f 2 \leftarrow$ IntervalFFT $(i 2)$ |  |
| $\operatorname{prod} \leftarrow \operatorname{Mul}(f 1, f 2)$ | $\triangleright$ point-wise multiplication |
| raw_retv $\leftarrow$ IntervalIFFT (prod) |  |
| $z \leftarrow \mathbf{C a r r y}$ (raw_retv) | $\triangle$ Propagate carries |

## Our Optimization in algorithm-level:

- IntervalFFT uses Decimation-in-frequency (DIF) and

IntervalIFFT uses Decimation-in-time (DIT) $\rightarrow$ avoid bit reversal operation

- Used Real FFT to help decrease the memory usage



## Our Optimization in software-level

- Used double-double (128 bits floating point data type) to provide higher precision
- Packed up more bits for an element to lower the FFT size
- Used shared memory for GPU to decrease global memory accesses
- Used OpenMP for CPU to parallelize


## Our Goal:

- Support multiplication of billions of digits input
- Achieve comparable or better performance in contrast with GMP



## Future work

- We plan to use pruned FFTs to further decrease our memory usage and speed up our algorithm
- We plan to combine our algorithm with Karatsuba to help roundoff errors
- We plan to use higher radix FFT in our algorithm

