Optimal Power Flow over Radial Networks

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Motivation

Semidefinite relaxation

- Bus injection model
- Conic relaxation
 - Branch flow model







Source: Rosa Yang, EPRI



OPF is solved routinely to determine

- How much power to generate where
- Market operation & pricing
- Parameter setting, e.g. taps, VARs
- Non-convex and hard to solve
 - Huge literature since 1962
 - In practice, operators often solve linearized model and verify using AC power flow model

Core of many problems
 OPF, LMP, Volt/VAR, DR, EV, planning ...



Wind power over land (exc. Antartica) 70 – 170 TW



Solar power over land 340 TW

Worldwide

energy demand: 16 TW

electricity demand: 2.2 TW

wind capacity (2009): 159 GW

grid-tied PV capacity (2009): 21 GW

Source: Renewable Energy Global Status Report, 2010 Source: M. Jacobson, 2011



Current control paradigm works well today

- Low uncertainty, few active assets to control
- Centralized, open-loop, human-in-loop, worst-case preventive
- Schedule supplies to match loads

Future needs

- Fast computation to cope with rapid, random, large fluctuations in supply, demand, voltage, freq
- Simple algorithms to scale to large networks of active DER
- Real-time data for adaptive control



Must close the loop

- Real-time feedback control, risk-limiting
- Driven by uncertainty of renewables

Must be scalable

- Distributed & decentralized optimization
- Orders of magnitude more endpoints that can generate, compute, communicate, actuate

Control and optimization framework

- Theoretical foundation for a holistic framework that integrates engineering + economics
- Systematic algorithm design, understandable global behavior
- Clarify ideas, explore structures, suggest direction



Motivation

Static capacitor control cannot cope with rapid random fluctuations of PVs on distr circuits

Inverter control

- Much faster & more frequent
- IEEE 1547 does not optimize VAR currently (unity PF)







Load and Solar Variation





Empirical distribution of (load, solar) for Calabash



Implication: reduced likelihood of violating voltage limits or VAR flow constraints

Energy savings



Summary

RESULTS FOR SOME VOLTAGE TOLERANCE THRESHOLDS

Voltage Drop	Annual Hours Saved Spending	Average Power
Tolerance(pu)	Outside Feasibility Region	Saving (%)
0.03	842.9	3.93%
0.04	160.7	3.67%
0.05	14.5	3.62%

- More reliable operation
- Energy savings





Radial networks computationally simple

- Exploit tree graph & convex relaxation
- Real-time scalable control promising



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Problem formulation

Carpentier 1962

Computational techniques:

- Dommel & Tinney 1968
- Surveys: Huneault et al 1991, Momoh et al 2001, Pandya et al 2008

SDP formulation (bus injection model):

- Bai et al 2008, 2009, Lavaei et al 2010
- Bose et al 2011, Sojoudi et al 2011, Zhang et al 2011
- Lesieutre et al 2011

Branch flow model

Baran & Wu 1989, Chiang & Baran 1990, Farivar et al 2011



Models: Kirchoff's law





Nodes *i* and *j* are linked with an admittance $y_{ij} = g_{ij} - \mathbf{i}b_{ij}$

$$Y_{ij} = \begin{cases} y_{ii} + \sum_{j \sim i} y_{ij}, & \text{if } i = j \\ -y_{ij}, & \text{if } i \sim j \\ 0 & \text{otherwise} \end{cases}$$

• Kirchhoff's Law: I = YV



 $\sum f_k(P_k^g) \leftarrow \text{Generation cost}$ min $k \in G$ $\left(P_k^g, Q_k^g, V_k\right)$ over subject to $\underline{P}_k^g \leq \underline{P}_k^g \leq \overline{P}_k^g$ Generation power constraints $Q_k^g \leq Q_k^g \leq \overline{Q}_k^g$ $\underline{V}_k \leq |V_k| \leq V_k$ Voltage magnitude constraints KVL/KCL power balance — nonconvexity



In terms of V:

$$P_{k} = \operatorname{tr} \Phi_{k} V V^{*}$$
$$Q_{k} = \operatorname{tr} \Psi_{k} V V^{*}$$

$$\Phi_k := \left(\frac{Y_k^* + Y_k}{2}\right)$$
$$\Psi_k := \left(\frac{Y_k^* - Y_k}{2\mathbf{i}}\right)$$

$$\min \sum_{k \in G} \operatorname{tr} M_{k}VV^{*}$$
over V
s.t.
$$\underline{P}_{k}^{g} - P_{k}^{d} \leq \operatorname{tr} \Phi_{k}VV^{*} \leq \overline{P}_{k}^{g} - P_{k}^{d}$$

$$\underline{Q}_{k}^{g} - Q_{k}^{d} \leq \operatorname{tr} \Psi_{k}VV^{*} \leq \overline{Q}_{k}^{g} - Q_{k}^{d}$$

$$\underline{V}_{k}^{2} \leq \operatorname{tr} J_{k}VV^{*} \leq \overline{V}_{k}^{2}$$

Key observation [Bai et al 2008]: OPF = rank constrained SDP



$$\min \sum_{k \in G} \operatorname{tr} M_k W$$

over *W* positive semidefinite matrix

s.t.
$$\underline{P}_{k} \leq \operatorname{tr} \Phi_{k} W \leq \overline{P}_{k}$$
$$\underline{Q}_{k} \leq \operatorname{tr} \Psi_{k} W \leq \overline{Q}_{k}$$
$$\underline{V}_{k}^{2} \leq \operatorname{tr} J_{k} W \leq \overline{V}_{k}^{2}$$
$$W \geq 0, \quad \operatorname{rank} W \equiv 1$$

convex relaxation: SDP





$$\min \sum_{k \in G} \operatorname{tr} M_k W$$

over *W* positive semidefinite matrix

s.t.
$$\underline{P}_{k} \leq \operatorname{tr} \Phi_{k} W \leq \overline{P}_{k}$$

 $\underline{Q}_{k} \leq \operatorname{tr} \Psi_{k} W \leq \overline{Q}_{k}$
 $\underline{V}_{k}^{2} \leq \operatorname{tr} J_{k} W \leq \overline{V}_{k}^{2}$
 $W \geq 0$
 $\underline{P}_{k} \leq \operatorname{tr} J_{k} W \leq \overline{V}_{k}^{2}$
 $\underline{\gamma}_{k}, \overline{\gamma}_{k}$
Lagrange
multipliers

$$A(\lambda_k, \mu_k, \gamma_k) := \sum_{k \in G} M_k + \sum_k (\lambda_k \Phi_k + \mu_k \Psi_k + \gamma_k J_k)$$



Theorem

If A^{opt} has rank n-1 then
W^{opt} has rank 1, SDP relaxation is exact
Duality gap is zero
A globally optimal V^{opt} can be recovered

IEEE test systems (essentially) satisfy the condition!

J. Lavaei and S. H. Low: Zero duality gap in optimal power flow problem. Allerton 2010, TPS 2011



Suppose

□ tree (radial) network

no lower bounds on power injections

Theorem

- A^{opt} always has rank n-1
- W^{opt} always has rank 1 (exact relaxation)
- □ OPF always has z₽red duality gap

Globally optimal solvable efficiently

S. Bose, D. Gayme, S. H. Low and M. Chandy, OPF over tree networks. Allerton 2011



Suppose

□ tree (radial) network

no lower bounds on power injections

Theorem

- A^{opt} always has rank n-1
- W^{opt} always has rank 1 (exact relaxation)

□ OPF always has z₽red duality gap

Globally optimal solvable efficiently

Also: B. Zhang and D. Tse, Allerton 2011 S. Sojoudi and J. Lavaei, submitted 2011





graph of QCQP $G(C,C_k)$ has edge $(i,j) \Leftrightarrow$

 $C_{ij} \neq 0$ or $[C_k]_{ij} \neq 0$ for some k

QCQP over tree $G(C, C_k)$ is a tree





Semidefinite relaxation
minmintr CWover $W \ge 0$ s. t.tr $C_k W \le b_k$ $k \in K$





Key assumption $(i, j) \in G(C, C_k): 0 \notin \text{ int conv hull } (C_{ij}, [C_k]_{ij}, \forall k)$

Theorem

Semidefinite relaxation is exact for QCQP over tree S. Bose, D. Gay

S. Bose, D. Gayme, S. H. Low and M. Chandy, submitted March 2012





Theorem

- A^{opt} always has rank *n*-1
 □ W^{opt} always has rank 1 (exact relaxation)
- □ OPF always has z∯rot duality gap
- Clabally antimal calvable afficiently





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Motivation

Semidefinite relaxationBus injection model

Conic relaxation

Branch flow model







Kirchoff's Law:
$$S_{ij} = \sum_{k:j \sim k} S_{jk} + z_{ij} |I_{ij}|^2 + s_j^c - s_j^g$$

line load - gen loss



Kirchoff's Law:
$$S_{ij} = \sum_{k:j \sim k} S_{jk} + z_{ij} |I_{ij}|^2 + S_j^c - S_j^g$$

Ohm's Law:
$$V_j = V_i - z_{ij}I_{ij}$$
 $S_{ij} = V_iI_{ij}^*$



OPF using branch flow model

$$\min \sum_{i \sim j} r_{ij} l_{ij} + \sum_{i} \alpha_{i} v_{i}$$

over (S, I, V, s^{g}, s^{c})
s. t. $\underline{s}_{i}^{g} \leq s_{i}^{g} \leq \overline{s}_{i}^{g}$ $\underline{s}_{i} \leq s_{i}^{c}$
 $\underline{v}_{i} \leq v_{i} \leq \overline{v}_{i}$

Kirchoff's Law:
$$S_{ij} = \sum_{k:j \sim k} S_{jk} + z_{ij} |I_{ij}|^2 + S_j^c - S_j^g$$

Ohm's Law:
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1. Angle relaxation

Angles of I_i , V_i eliminated ! Points relaxed to circles

demands $P_{ii} = \sum P_{ik} + r_{ii} |I_{ii}|^2 + p_i^c - p_i^g$ $k:i \sim k$ $Q_{ii} = \sum Q_{ik} + x_{ii} |I_{ii}|^2 + q_i^c - q_i^g$ *k*:*i*~*k* $|V_{i}|^{2} = |V_{i}|^{2} + 2(r_{ii}P_{ii} + x_{ii}Q_{ii}) - (r_{ii}^{2} + x_{ii}^{2})|I_{ii}|^{2}$ $|I_{ij}|^2 = \left(\frac{P_{ij}^2 + Q_{ij}^2}{|V_{.}|^2}\right)$ Baran and Wu 1989 for radial networks

1. Angle relaxation

Angles of I_i , V_i eliminated ! Points relaxed to circles

$$P_{ij} = \sum_{k:j \sim k} P_{jk} + r_{ij} \left| I_{ij} \right|^2 + p_j^c - p_j^g \qquad \text{generation}$$

$$Q_{ij} = \sum_{k:j \sim k} Q_{jk} + x_{ij} \left| I_{ij} \right|^2 + q_j^c - q_j^g \qquad \text{VAR contro}$$

$$|V_i|^2 = \left| V_j \right|^2 + 2\left(r_{ij} P_{ij} + x_{ij} Q_{ij} \right) - \left(r_{ij}^2 + x_{ij}^2 \right) \left| I_{ij} \right|^2$$

$$|I_{ij}|^2 = \left(\frac{P_{ij}^2 + Q_{ij}^2}{\left| V_i \right|^2} \right) \qquad \underbrace{\underline{S}_i^g} \leq \underline{S}_i^g \leq \overline{S}_i^g \qquad \underline{S}_i \leq \underline{S}_i^c$$

1. Angle relaxation



$$\min \sum_{i \sim j} r_{ij} l_{ij} + \sum_{i} \alpha_{i} v_{i}$$

over (S, l, v, s^{g}, s^{c})
s. t. $P_{ij} = \sum_{k:j \sim k} P_{jk} + r_{ij} l_{ij} + p_{j}^{c} - p_{j}^{g}$
 $Q_{ij} = \sum_{k:j \sim k} Q_{jk} + x_{ij} l_{ij} + q_{j}^{c} - q_{j}^{g}$

- Linear objective
- Linear constraints
- Quadratic equality

$$v_{i} = v_{j} + 2\left(r_{ij}P_{ij} + x_{ij}Q_{ij}\right) - \left(r_{ij}^{2} + x_{ij}^{2}\right)l_{ij}$$

$$l_{ij} = \left(\frac{P_{ij}^{2} + Q_{ij}^{2}}{v_{i}}\right), \qquad \underline{s}_{i} \leq s_{i}^{g} \leq \overline{s}_{i}$$

$$\underline{v}_{i} \leq v_{i} \leq \overline{v}_{i}, \qquad \underline{s}_{i} \leq \overline{s}_{i}^{c}$$

2. SOCP relaxation

$$\min \sum_{i \sim j} r_{ij} l_{ij} + \sum_{i} \alpha_{i} v_{i}$$

$$over \quad (S, l, v, s^{g}, s^{c})$$

$$s. t. \quad P_{ij} = \sum_{k:j \sim k} P_{jk} + r_{ij} l_{ij} + p_{j}^{c} - p_{j}^{g}$$

$$Q_{ij} = \sum_{k:j \sim k} Q_{jk} + x_{ij} l_{ij} + q_{j}^{c} - q_{j}^{g}$$

$$v_{i} = v_{j} + 2 \left(r_{ij} P_{ij} + x_{ij} Q_{ij} \right) - \left(r_{ij}^{2} + x_{ij}^{2} \right) l_{ij}$$

$$lity \quad l_{ij} \ge \left(\frac{P_{ij}^{2} + Q_{ij}^{2}}{v_{i}} \right) \qquad \underline{s}_{i} \le s_{i}^{g} \le \overline{s}_{i}$$

$$\underline{v}_{i} \le v_{i} \le \overline{v}_{i}, \qquad \underline{s}_{i} \le s_{i}^{c}$$

Quadratic inequality

OPF over radial networks

Theorem

Both relaxation steps are exact

- SOCP relaxation is (convex and) exact
- Phase angles can be uniquely determined

Original OPF problem has zero duality gap

M. Farivar, C. Clarke, S. H. Low and M. Chandy, Inverter VAR control for distribution systems with renewables. SmartGridComm 2011

What about mesh networks ??

M. Farivar and S. H. Low, submitted March 2012



OPF using branch flow model

$$\min \sum_{i \sim j} r_{ij} l_{ij} + \sum_{i} \alpha_{i} v_{i}$$

over (S, I, V, s^{g}, s^{c})
s. t. $\underline{s}_{i}^{g} \leq s_{i}^{g} \leq \overline{s}_{i}^{g}$ $\underline{s}_{i} \leq s_{i}^{c}$
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Kirchoff's Law:
$$S_{ij} = \sum_{k:j \sim k} S_{jk} + z_{ij} |I_{ij}|^2 + S_j^c - S_j^g$$

Ohm's Law:
$$V_j = V_i - z_{ij}I_{ij}$$
 $S_{ij} = V_iI_{ij}^*$



PF
$$\min_{x,s} f(\hat{h}(x))$$
 s.t. $x \in \mathbf{X}, s \in \mathbf{S}$

OPF-ar
$$\min_{x,s} f(\hat{h}(x))$$
 s.t. $x \in \mathbf{Y}, s \in \mathbf{S}$

OPF-ps
$$\min_{x,s,\phi} f(\hat{h}(x))$$
 s.t. $x \in \overline{\mathbf{X}}, s \in \mathbf{S}$

Theorem

• $\overline{\mathbf{X}} = \mathbf{Y}$

 \bigcap

 Need phase shifters only outside spanning tree





OPF
$$\min_{x,s} f(\hat{h}(x))$$
 s.t. $x \in \mathbf{X}$

PF-ar
$$\min_{x,s} f(\hat{h}(x))$$
 s.t. $x \in \mathbf{Y}$

OPF-ps
$$\min_{x,s,\phi} f(\hat{h}(x))$$
 s.t. $x \in \overline{\mathbf{X}}$

Theorem

- $\overline{\mathbf{X}} = \mathbf{Y}$
- Need phase shifters only outside spanning tree





Radial networks computationally simple

- Exploit tree graph & convex relaxation
- Real-time scalable control promising

Mesh networks can be convexified

- Design for simplicity
- Need few phase shifters (sparse topology)