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## Vector Quantization and Subband Coding



18-796 Multimedia Communications:  
Coding, Systems, and Networking

*Prof. Tsuhan Chen*  
tsuhan@ece.cmu.edu

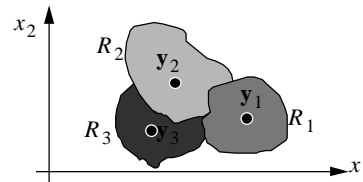
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## Vector Quantization



## Vector Quantization (VQ)

- Each image block has  $N$  pels
- Consider each image block as an  $N$ -D vector  $\mathbf{x}$



- Quantization:  $\mathbf{x} \rightarrow \mathbf{y}_k$  if  $\mathbf{x} \in R_k$
- $\mathbf{y}_k$ : codewords or code vectors
- The set of  $\mathbf{y}_k$  is called a codebook

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## Rate and Distortion

- If the number of codewords is  $K$ , then the number of bits required to send one vector is  $\log_2 K$
- Rate  $R$ 
  - $R$  bits per pixel
  - $NR$  bits for one vector, so  $\log_2 K = NR$ , i.e.,  $\log K = 2^{NR}$
- Distortion  $D$ 
  - Given the probability density function  $p(\mathbf{x})$  and distortion measure  $d(\mathbf{x}, \mathbf{y})$ , the average distortion is

$$D = \sum_{k=1}^K \int_{R_k} d(\mathbf{x}, \mathbf{y}_k) p(\mathbf{x}) d\mathbf{x}$$

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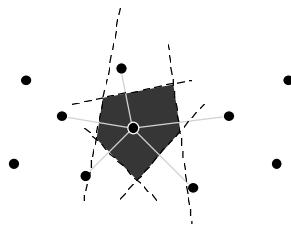
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- Given  $\mathbf{y}_k$ ,  $R_k$  should be chosen such that

$$R_k = \left\{ \mathbf{x}: d(\mathbf{x}, \mathbf{y}_k) \leq d(\mathbf{x}, \mathbf{y}_j) \forall j \neq k \right\}$$

= the set of  $\mathbf{x}$  for which  $\mathbf{y}_k$  is the nearest point

- For  $L_2$  norm, i.e.,  $d(\mathbf{x}, \mathbf{y}) = \frac{1}{N} \sum_{i=1}^N (x_i - y_i)^2$  we get



**Convex Polytope  
(Voronoi Cell)**

Q: How about  $L_1$  or  $L_\infty$ ?

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- Given  $R_k$ ,  $\mathbf{y}_k$  should be chosen such that

$$\int_{R_k} d(\mathbf{x}, \mathbf{y}_k) p(\mathbf{x}) d\mathbf{x} \text{ is minimum}$$

- With  $L_2$  norm, we get

$$\mathbf{y}_k = \text{centroid of } R_k = \int_{R_k} \mathbf{x} p(\mathbf{x}) d\mathbf{x}$$

- In the discrete case, optimal  $\mathbf{y}_k$  is the average of the vectors in  $R_k$

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## Generalized Lloyd Algorithm (LBG Algorithm, K-means Algorithm)

- Linde, Buzo, and Gray, 1980
- Given  $p(\mathbf{x})$ , or given a set of training vectors
  - (1) Start with an initial set of  $\mathbf{y}_k$ , i.e., initial codebook
  - (2) With the current  $\mathbf{y}_k$ , calculate the region  $R_k$
  - (3) Replace each  $\mathbf{y}_k$  with the centroid of  $R_k$
  - (4) If the overall distortion  $D$  is lower than a threshold, stop.  
Otherwise, go to (2)
- Only gives local optimum. Proper choice of initial codebook is important

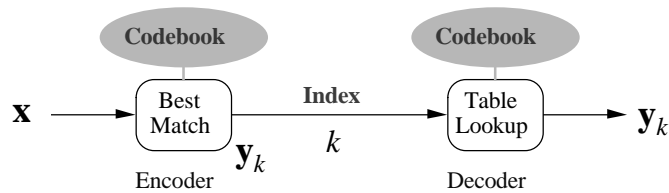
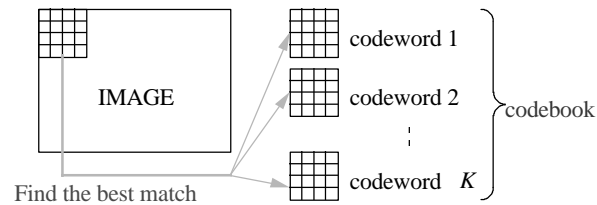
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- Choice of initial codebook
  - A representative subset of the training vectors
  - Scalar quantization in each dimension
  - Splitting...
- Nearest Neighbor (NN) algorithm [Equitz, 1984]
  - Start with the entire training set
  - Merge the two vectors that are closest into one vector equal to their mean
  - Repeat until the desired number of vectors is reached, or the distortion exceeds a certain threshold

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## Image Coding



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## Properties of VQ

- Codebook design is very complex
  - 4x4 blocks at 1 bpp:  $2^{16}$  codewords
  - 16 images of size 256x256:  $2^{16}$  training vectors (4x4 each)
  - Codebook size:  $2^{16} \times 4 \times 4 \times 8 \text{ bits} = 8.3 \text{ Mbits}$
- More useful for low bitrate
  - 4x4 blocks at 0.5 bpp:  $2^8 = 256$  codewords
  - One 256x256 image: 4096 training vectors
  - Codebook size:  $256 \times 4 \times 4 \times 8 \text{ bits} = 32.8 \text{ Kbits}$
- Simple decoder, complex encoder
  - Very good for image retrieval
- Poor performance on images not in the training set vs. overhead of sending the codebook

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## VQ Variants and Improvements

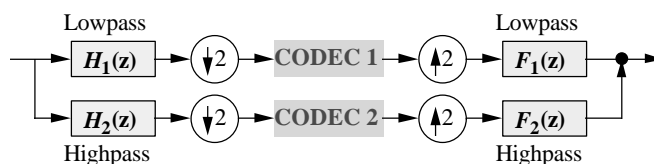
- Multistage VQ
- Product Codes
  - Send mean and variance separately
- Classified VQ
  - Edges, texture areas, flat areas
- Predictive VQ
- VQ for color images
  - Exploit correlation among color components, e.g. R,G,B
  - YUV components are practically uncorrelated

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## Subband Coding



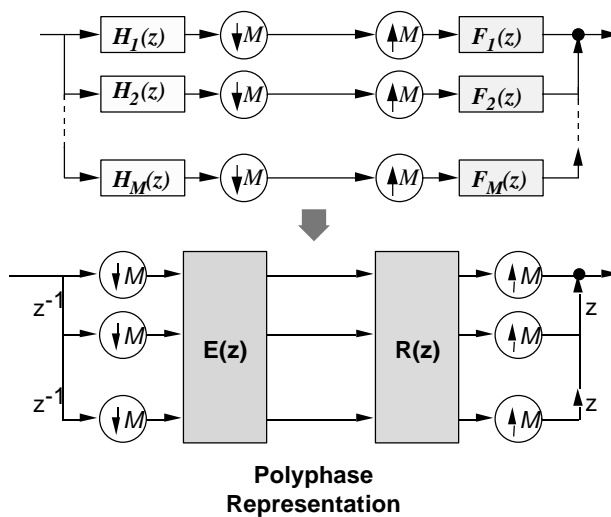
## Subband Coding



- Decompose the signal in the frequency domain
- Critical downsampling (maximal decimation) maintains the number of samples in the subbands
- Wavelet coding: Recursively apply subband decomposition to the low freq band
- 2-D: Separable filtering to get 4 bands: LL, LH, HL, HH

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## Subband Coding vs. Transform Coding



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- Perfect reconstruction (PR) is obtained when  $\mathbf{R}(z)=\mathbf{E}^{-1}(z)$
- When  $\mathbf{E}(z)$  and  $\mathbf{R}(z)$  are constant matrices, subband coding degenerated to blocked-based operation, i.e., transform coding
- In particular, if  $\mathbf{E}(z)$  is a DCT matrix and  $\mathbf{R}(z)$  is IDCT, this becomes DCT coding
- Subband coding can be viewed as transform coding with overlapped blocks. So, it can exploit correlation of pixels at longer range
- Coding Artifacts:
  - Transform Coding: blocking
  - Subband Coding: ringing, contouring

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## Optimal Bit Allocation

- We can allocate different bit rates to the subbands based on their properties
- Assume that we apply scalar quantization with bitrate  $b_k$  to the subbands  $x_k$ , then the quantization error is

$$s_{q_k}^2 = c \times 2^{-2b_k} s_{x_k}^2$$

- The overall quantization error is  $s_q^2 = \frac{1}{M} \sum_{k=1}^M s_{q_k}^2$
- The overall bitrate is

$$b = \frac{1}{M} \sum_{k=1}^M b_k$$

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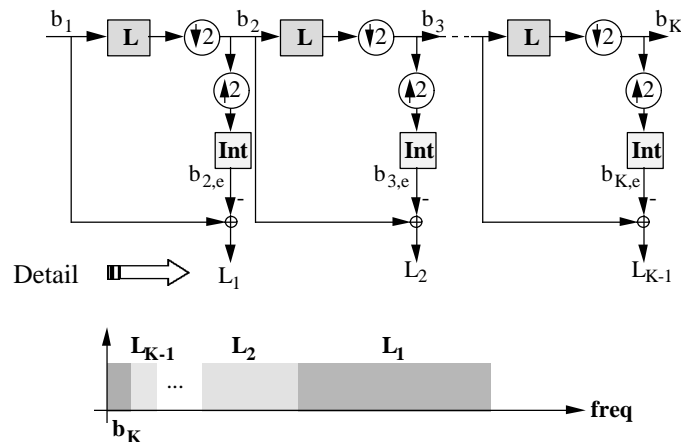
$$\begin{aligned}
 s_q^2 &\geq \left( \prod_{k=1}^M s_{q_k}^2 \right)^{1/M} \quad (\text{AM-GM inequality}) \\
 &= c \left( \prod_{k=1}^M 2^{-2b_k} s_{x_k}^2 \right)^{1/M} = c \left( 2^{-2 \sum b_k / M} \right) \left( \prod_{k=1}^M s_{x_k}^2 \right)^{1/M} \\
 &= c \times 2^{-2b} \left( \prod_{k=1}^M s_{x_k}^2 \right)^{1/M} \quad (\text{a constant for given signal and filter bank})
 \end{aligned}$$

- Equality holds if and only if  $s_{q_k}^2 = s_q^2 \quad \forall k$
- Optimal bit allocation  $b_k = \frac{1}{2} \log \frac{c \times s_{x_k}^2}{s_q^2}$
- Gain =  $\frac{\frac{1}{M} \sum_{k=1}^M s_{x_k}^2}{\left( \prod_{k=1}^M s_{x_k}^2 \right)^{1/M}} \geq 1$  No gain if  $s_{x_k}^2$  are identical

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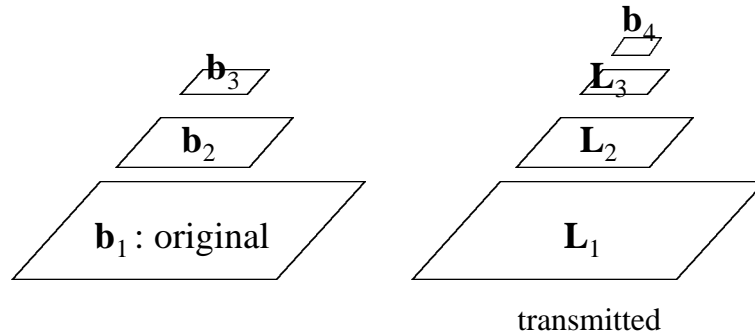
## Pyramid Coding



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## The “Pyramid”

- Consider the 2-D case



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- Number of samples is 33% more

$$N + \frac{N}{4} + \frac{N}{16} + \dots \approx \frac{4}{3}N$$

- Non-critical sampling
- PR is always possible
  - No matter how L and Int are designed
- Progressive transmission is possible

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  - John Woods, ed., *Subband Image Coding*, Kluwer Academic Publishers, 1991
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