All You Ever Wanted to Know About Dynamic Taint Analysis & Forward Symbolic Execution (but might have been afraid to ask)

(Yes, we were trying to overflow the title length field on the submission server)

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A ***Few Things*** You Need to Know About Dynamic Taint Analysis & Forward Symbolic Execution (but might have been afraid to ask)

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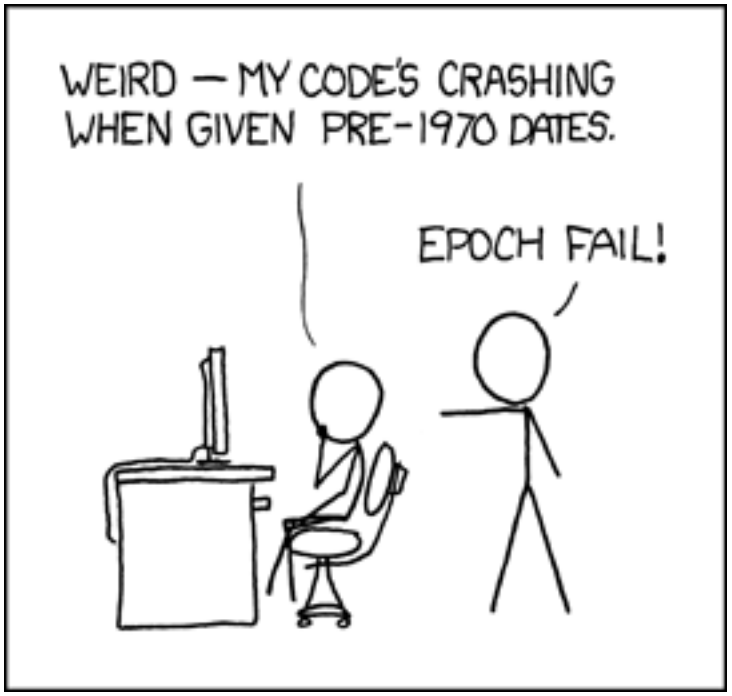
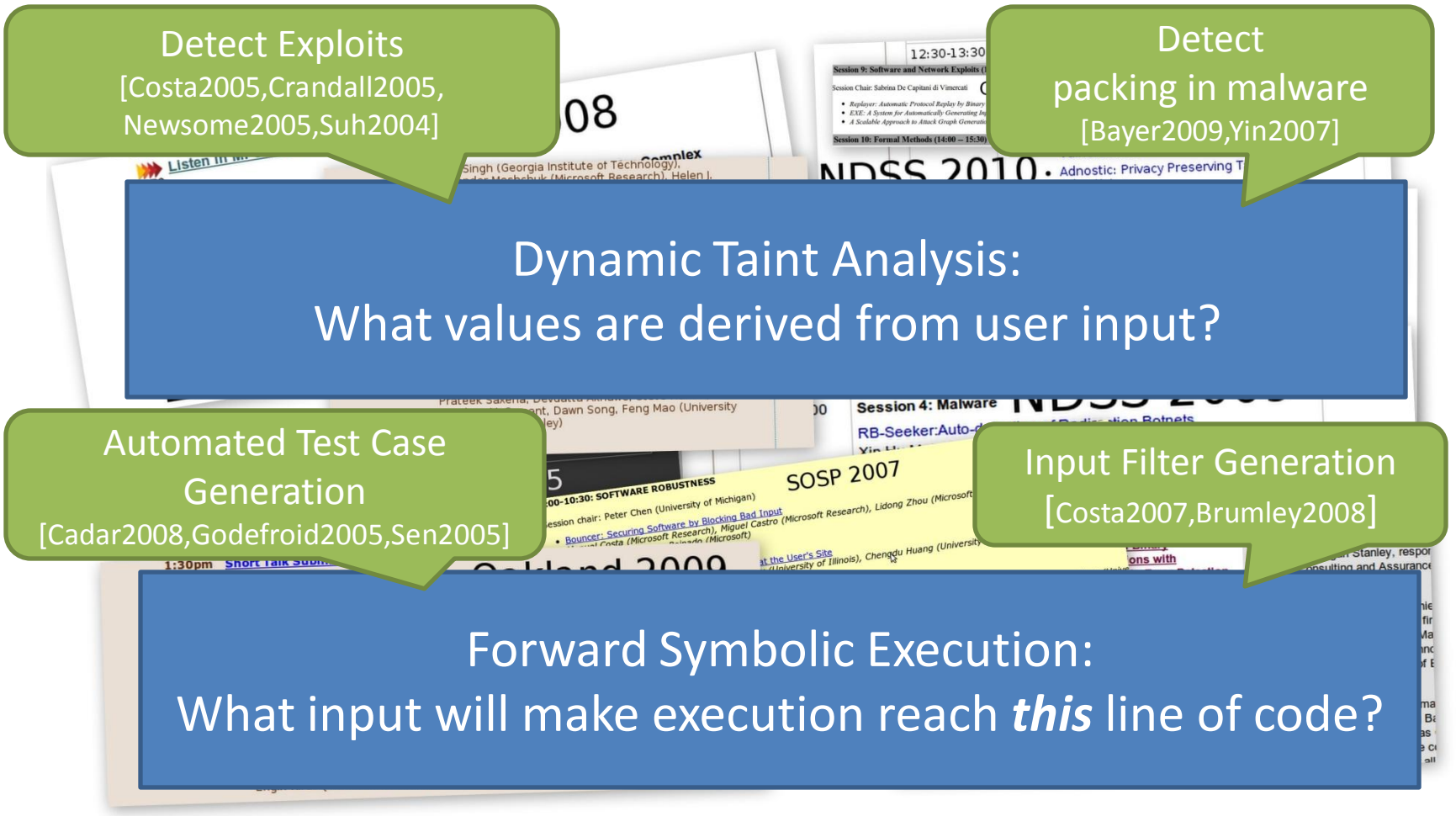
The Root of All Evil

Humans write programs



This Talk:
Computers Analyzing Programs Dynamically at Runtime

Two Essential Runtime Analyses



Our Contributions

Computers Analyzing Programs Dynamically at Runtime

Dynamic Taint Analysis:
Is this value affected by user input?

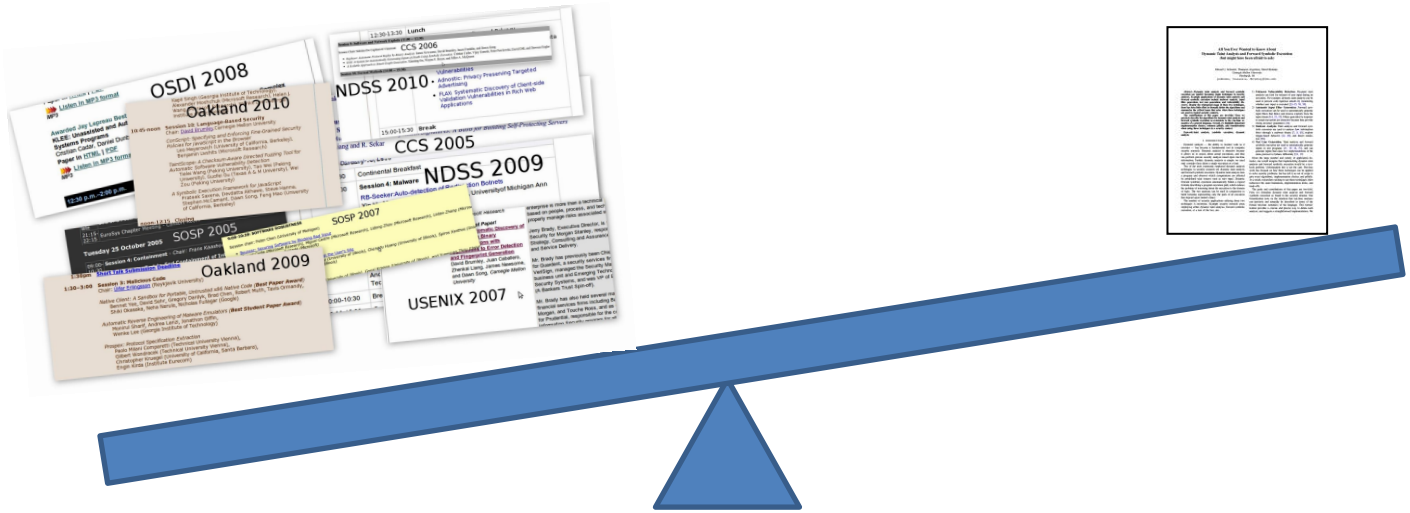
Forward Symbolic Execution:
What input will make execution reach ***this*** line of code?

1: Turn English descriptions into an ***algorithm***
   - Operational Semantics

2: Algorithm highlights caveats, issues, and unsolved problems that are deceptively hard

Our Contributions (cont'd)

3: Systematize recurring themes in a wealth of previous work



Dynamic Taint Analysis: What values are derived from user input?

1. How it works – example
2. Desired properties
3. Example issue. Paper has many more.

tainted untainted

x = get_input()

y = x + 42

…

goto y

Input is tainted

Δ

Var	Val
x	7

Taint Introduction

$$\text{Input}\ \frac{t = \text{IsUntrusted}(src)}{\text{get\_input}(src)\downarrow t}$$

τ

Var	Tainted?
x	T

tainted untainted

x = get_input( )

y = x + 42

…

goto y

Data derived from user input is tainted

Δ

Var	Val
x	7
y	49

Taint Propagation

$$\text{BinOp}\ \frac{t_1 = \tau[x_1]\ ,\ t_2 = \tau[x_2]}{x_1 + x_2 \downarrow t_1 \vee t_2}$$

τ

Var	Tainted?
x	T
y	T

tainted untainted

x = get_input()

y = x + 42

…

goto y

Policy Violation Detected

Taint Checking

$P_{goto}(t_a) = \neg t_a$

(Must be true to execute)

Δ

Var	Val
x	7
y	49

τ

Var	Tainted?
x	T
y	T

x = get_input( )

y = …

…

goto y

Jumping to overwritten return address

Different Use:
Exploit Detection

```
…
strcpy(buffer,argv[1]) ;
…
return ;
```

Memory Load

Variables

Δ

Var	Val
x	7

τ

Var	Tainted?
x	T

Memory

μ

Addr	Val
7	42

τ_μ

Addr	Tainted?
7	F

Problem: Memory Addresses

x = get_input( )

y = load(x)

…

goto y

All values derived from user input are tainted??

Δ

Var	Val
x	7

μ

Addr	Val
7	42

τ_μ

Addr	Tainted?
7	F

Policy 1: Taint depends only on the memory cell

x = get_ir

y = load(

…

goto y

Undertainting

Failing to identify tainted values

- e.g., missing exploits

Δ

Var	Val
x	7

Addr	Val
7	42

Taint Propagation

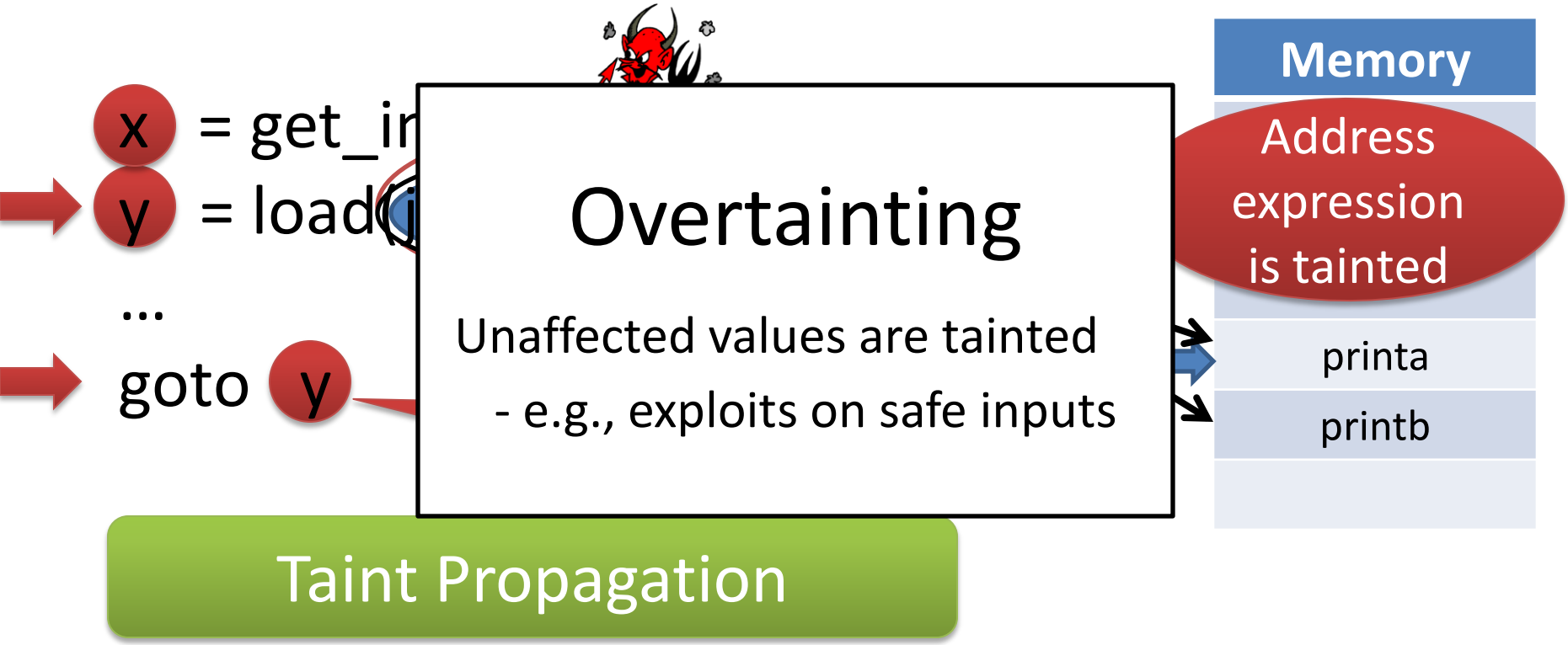
$$\text{Load}\ \frac{v = \Delta[x]\ ,\ t = \tau_\mu[v]}{\text{load}(x) \downarrow t}$$

τ_μ

Addr	Tainted?
7	F

Policy 2:

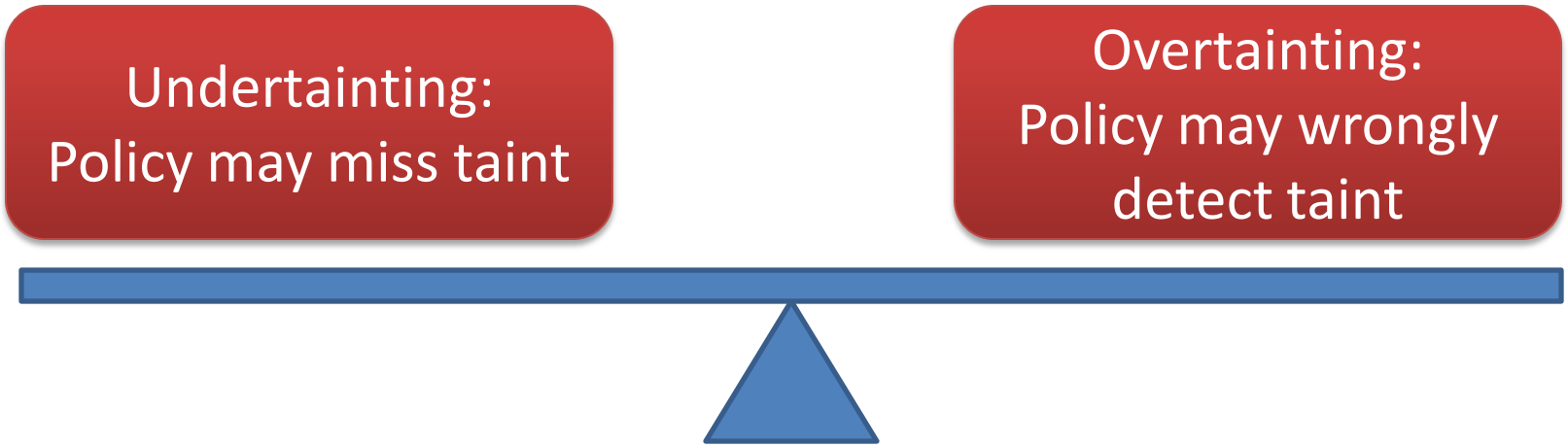
If either the address or the memory cell is tainted, then the value is tainted



Taint Propagation

$$\text{Load}\ \frac{v = \Delta[x]\ ,\ t = \tau_\mu[v],\ t_a = \tau[x]}{\text{load}(x) \downarrow t\ v\ t_a}$$

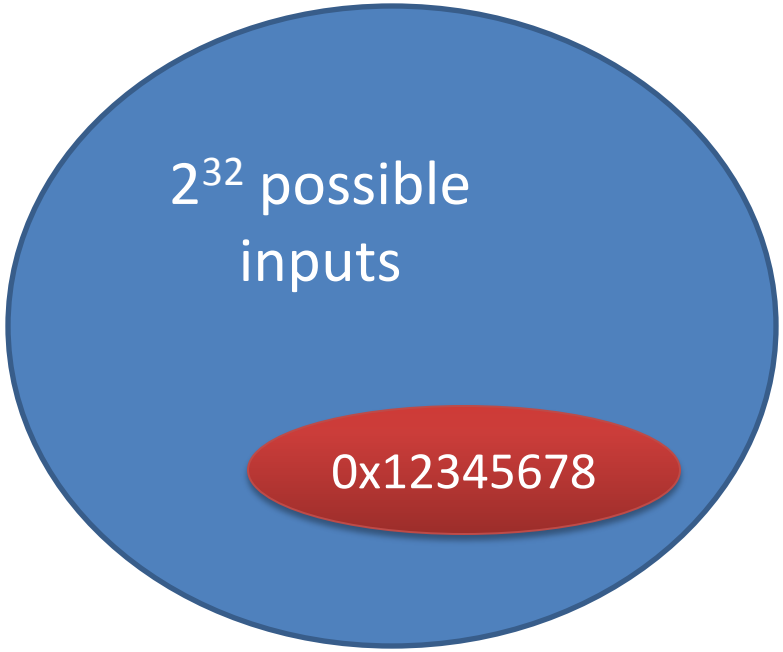
Research Challenge
State-of-the-Art is not perfect for all programs



Forward Symbolic Execution: What input will make execution reach ***this*** line of code?

- How it works – example
- Inherent problems of symbolic execution
- Proposed solutions

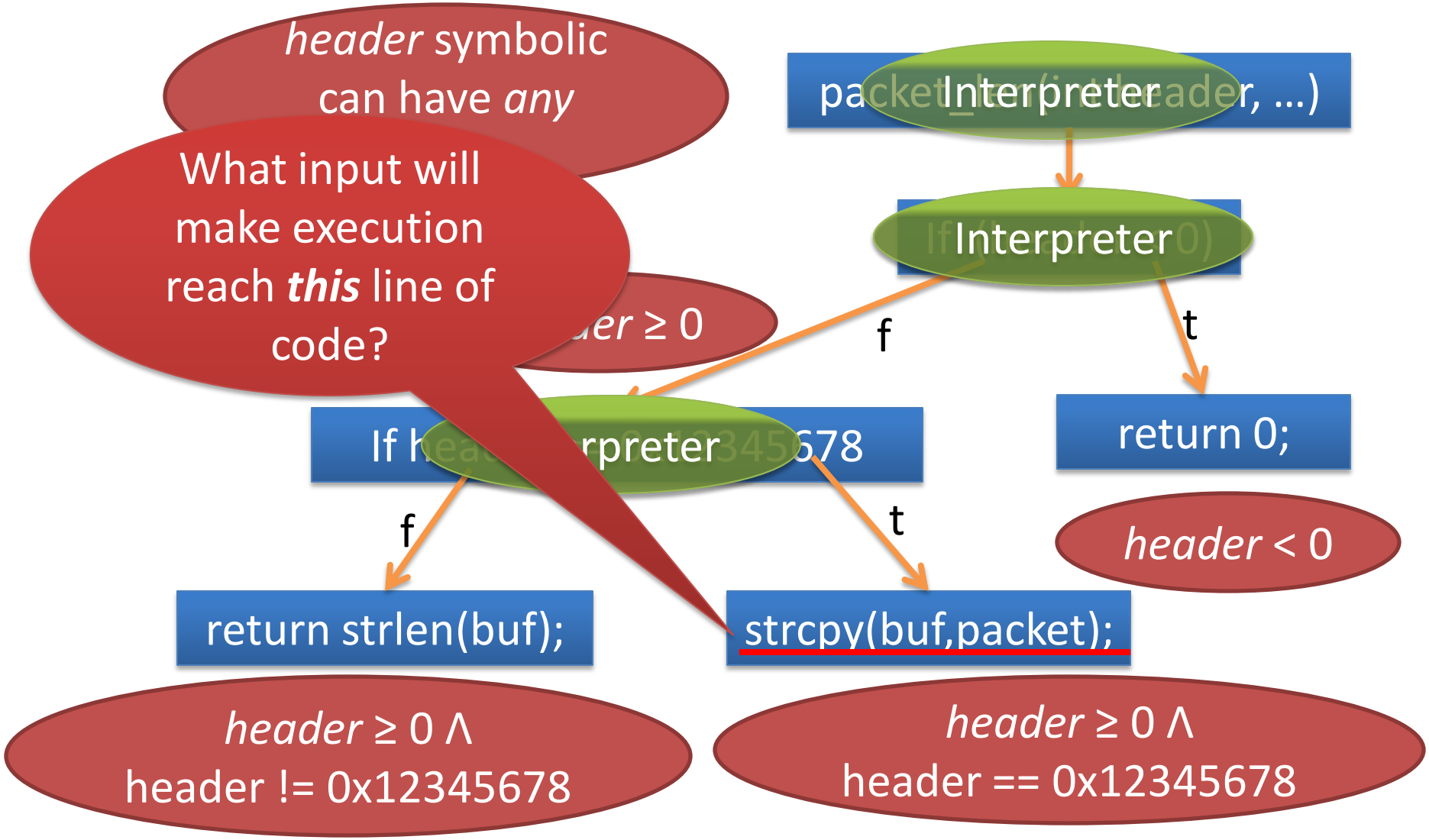
The Challenge



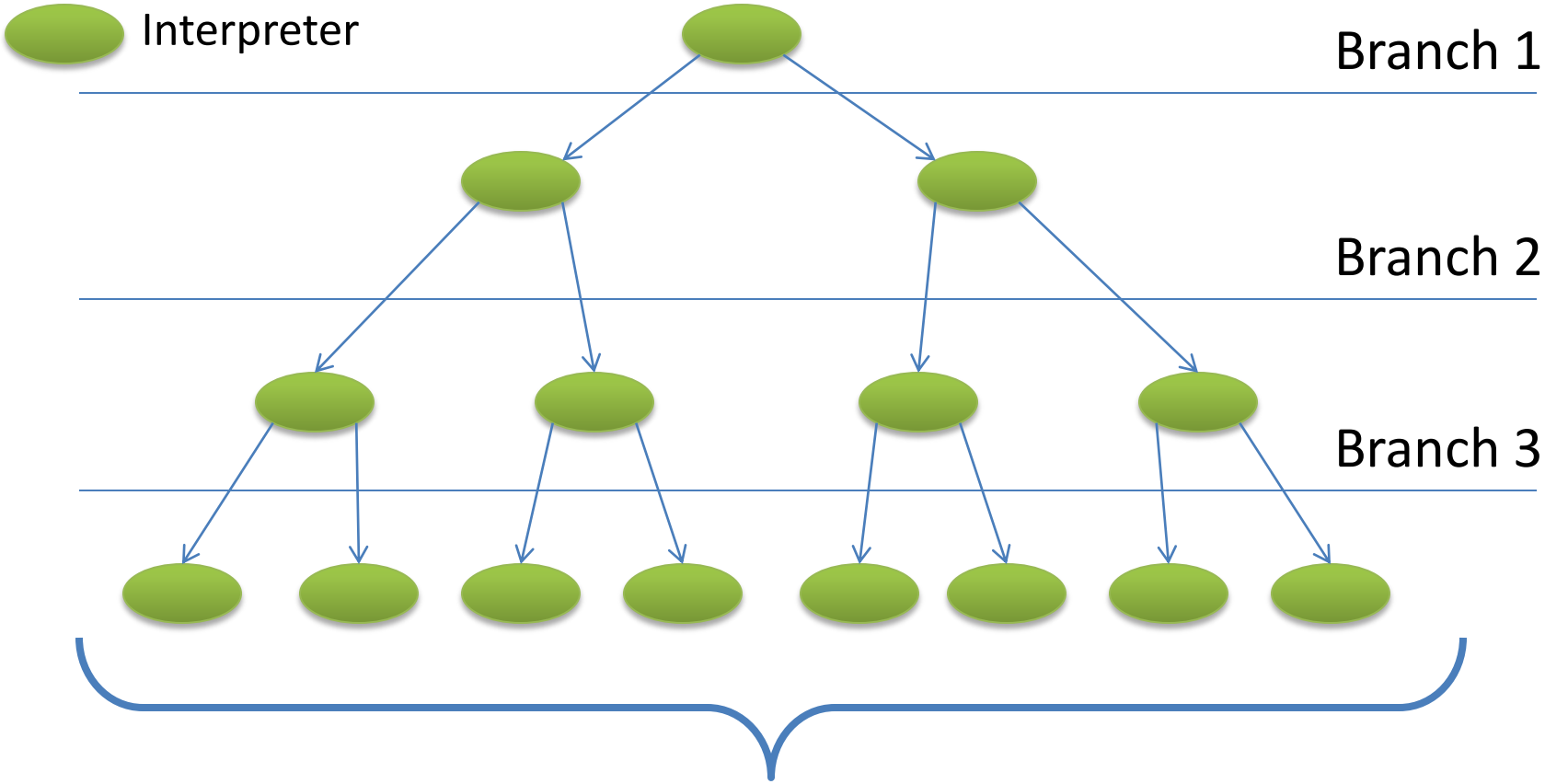

```
packet_len(int header, char *packet)
   char buf[2048]  = "…";
   if (header < 0)
            return 0;
   if (header == 0x12345678)
            strcpy(buf, packet);
   return strlen(buf);
```

Forward Symbolic Execution:
What input will make execution reach ***this*** line of code?

A Simple Example

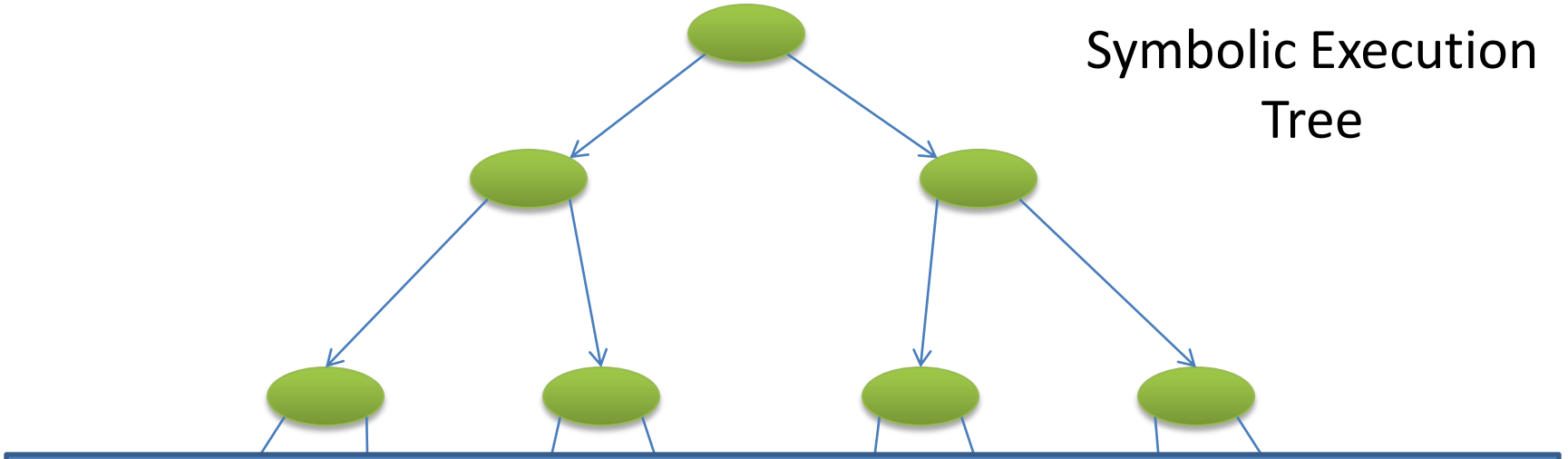


One Problem: Exponential Blowup Due to Branches



Exponential Number of Interpreters/formulas in # of branches

Path Selection Heuristics



However, these are heuristics. In the worst case all create an exponential number of formulas in the tree height.

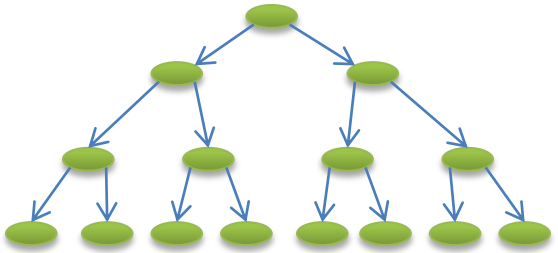
…

- Depth-First Search (bounded) ,Random Search [Cadar2008]
- Concolic Testing [Sen2005,Godefroid2008]

Symbolic Execution is *not* Easy

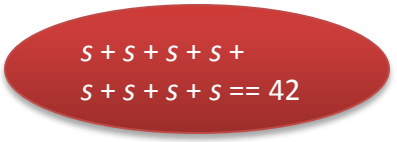
- Exponential number of interpreters/formulas

branching



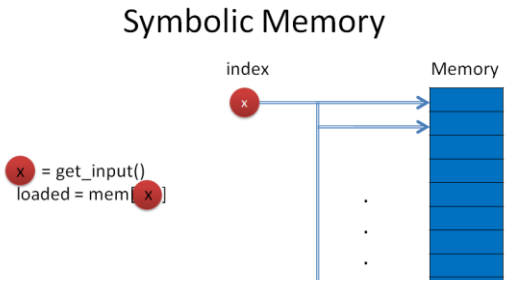
- Exponentially-sized formulas

substitution



- Solving a formula is NP-Complete!

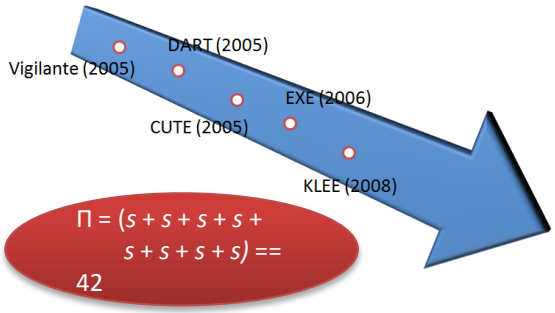
Other *Important* Issues



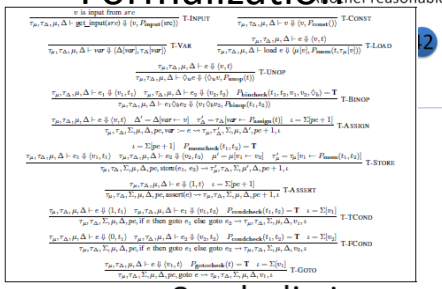
FORWARD SYMBOLIC EXECUTION

"Like normal execution, where inputs are substituted by symbolic variables"

King *et al.*, 1976

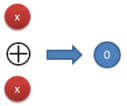


Formalization



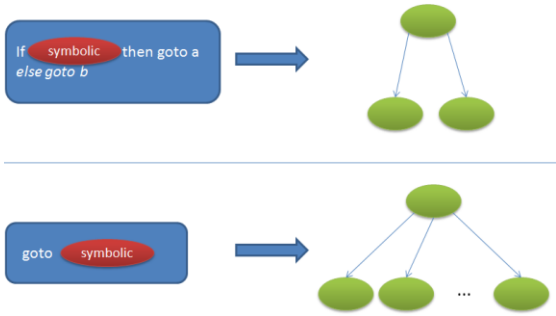
Sanitization

other reasonable assumption ?



More complex policies

Symbolic Jumps



Conclusion

- Dynamic taint analysis and forward symbolic execution used extensively in literature
  - Formal algorithm and what is done for each possible step of execution often not emphasized
- We provided a formal definition and summarized
  - Critical issues
  - State-of-the-art solutions
  - Common tradeoffs

Thank You!

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Questions?