

Algebraic Signal Processing Theory: 1-D Space

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Abstract—In our paper titled “Algebraic Signal Processing Theory: Foundation and 1-D Time” appearing in this issue of the IEEE TRANSACTIONS ON SIGNAL PROCESSING, we presented the algebraic signal processing theory, an axiomatic and general framework for linear signal processing. The basic concept in this theory is the signal model defined as the triple $(\mathcal{A}, \mathcal{M}, \Phi)$, where \mathcal{A} is a chosen algebra of filters, \mathcal{M} an associated \mathcal{A} -module of signals, and Φ is a generalization of the z -transform. Each signal model has its own associated set of basic SP concepts, including filtering, spectrum, and Fourier transform. Examples include infinite and finite discrete time where these notions take their well-known forms. In this paper, we use the algebraic theory to develop infinite and finite *space* signal models. These models are based on a *symmetric* space shift operator, which is distinct from the standard time shift. We present the *space* signal processing concepts of filtering or convolution, “ z -transform,” spectrum, and Fourier transform. For finite length *space* signals, we obtain 16 variants of space models, which have the 16 discrete cosine and sine transforms (DCTs/DSTs) as Fourier transforms. Using this novel derivation, we provide missing signal processing concepts associated with the DCTs/DSTs, establish them as precise analogs to the DFT, get deep insight into their origin, and enable the easy derivation of many of their properties including their fast algorithms.

Index Terms—Algebra, boundary condition, Chebyshev polynomials, convolution, discrete cosine transform (DCT), discrete sine transform (DST), Fourier transform, module, representation theory, shift, signal extension, signal model.

I. INTRODUCTION

STANDARD linear signal processing (SP) considers signals indexed by time (discrete or continuous) and time-invariant systems or filters. Associated with SP is the time shift operator, abstractly defined (in discrete form) as

$$q \diamond t_n = t_{n+1}. \quad (1)$$

The formulas for linear convolution and the discrete-time Fourier transform for infinite-length signals or for circular convolution and the discrete Fourier transform (DFT) for finite-length signals can be derived from this definition of the shift.

In this paper we show that an alternative linear SP framework can be derived from a different definition of the shift operator. This shift operates *undirected* or *symmetrically* in contrast to the directed operation of the time shift in (1). For this reason, we call it the *space shift*; it is abstractly defined as

$$q \diamond t_n = \frac{1}{2}(t_{n-1} + t_{n+1}). \quad (2)$$

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Accordingly, we derive for infinite- and finite-length signals the appropriate space SP notions including filtering or convolution, “ z -transforms,” spectrum, Fourier transforms, frequency response, and others. In the finite case, we explain the need for boundary conditions and identify 16 “natural” choices that have the 16 discrete cosine and sine transforms (DCTs/DSTs) as Fourier transforms. This establishes the DCTs/DSTs as exact analogs of the DFT, a satisfying alternative to the original derivation of the DCTs and DSTs as approximations to the Karhunen–Loève transform of a stationary process [2], [3]. The complete set of DCTs/DSTs was defined in [4] without derivation or motivation. In this paper, we jointly refer to the DCTs and DST as discrete trigonometric transforms (DTTs) even though this class is actually larger (e.g., it contains the real DFT and discrete Hartley transform).

We note that in other areas such as dynamic systems, it is common to consider different notions of shift [5].

We develop *space* SP as an instantiation of the algebraic signal processing theory (ASP), a general and axiomatic theory of (linear) SP presented in [1] and [6]. The central object in ASP is the signal model, defined as a triple $(\mathcal{A}, \mathcal{M}, \Phi)$, where \mathcal{A} is the filter space (an algebra), \mathcal{M} the signal space (an \mathcal{A} -module), and Φ generalizes the concept of z -transform. Many signal models are in principle possible, each with its own SP notions, including filtering, spectrum, or Fourier transform. ASP establishes that for finite signals and shift-invariant models, \mathcal{A} and \mathcal{M} are polynomial algebras $\mathbb{C}[x]/p(x)$, i.e., spaces of polynomials with multiplication modulo a fixed polynomial. For example, for the finite time model, which has the DFT as Fourier transform, both take the form $\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/(x^n - 1)$.

In [1], we explained how to derive signal models from a definition of the shift. Application to the time shift (1) yielded the well-known infinite and finite time signal models. In this paper, we derive signal models from the space shift (2). We identify and define the C -transform as the appropriate “ z -transform” and, for finite space signals, show that the 16 DTTs are the appropriate space Fourier transforms. As expected, the finite space signal models underlying the DTTs are again built from polynomial algebras. One application of the ASP interpretation of the DTTs is the easy derivation of many of their properties and their fast algorithms [7]–[9].

The DCT, type 3, was related to a polynomial algebra in [10]; all DTTs of types 1–4 were related to polynomial algebras in [11]; see also [12]. In all cases, no connection to signal processing was established.

Organization: We start with a brief overview of ASP in Section II. The focus will be on finite shift-invariant signal models that are built from polynomial algebras. In Sections III and IV, we derive the infinite and finite space models. The finite case is worked out in greater detail since it provides the underpinning of the frequently used DTTs including many of

their properties. An important variant of the DTTs, and their underlying signal models, is derived in Section V. Finally, we offer conclusions in Section VII.

II. ALGEBRAIC SIGNAL PROCESSING THEORY

We introduce the necessary background on the ASP and show infinite and finite time signal processing as examples. For a complete and detailed introduction we refer the reader to [1], [6]. For brevity we will denote linear signal processing by SP.

A. Signal Model

Algebra (Filter Space): An algebra \mathcal{A} is a vector space that is also a ring, i.e., it permits multiplication of elements and the distributivity law holds. Examples include the sets \mathbb{C}, \mathbb{R} of complex or real numbers and the set of polynomials with complex coefficients $\mathbb{C}[x]$. In SP, the set of filters is commonly assumed to be an algebra, with the multiplication being the concatenation of filters. We denote elements of algebras with h , the common symbol for filters in SP.

Module (Signal Space): Given an algebra \mathcal{A} , an \mathcal{A} -module \mathcal{M} is a vector space that permits an operation “ \cdot ” of \mathcal{A} on \mathcal{M} :

$$h \cdot s \in \mathcal{M}, \quad \text{for } h \in \mathcal{A}, \quad s \in \mathcal{M}. \quad (3)$$

Further, several properties such as the distributivity law have to hold [13]. In SP, the signal space is commonly assumed to be an \mathcal{A} -module, where \mathcal{A} is the associated space of filters. The operation \cdot denotes filtering; (3) ensures that filtering a signal $s \in \mathcal{M}$ with a filter $h \in \mathcal{A}$ yields again a signal.

A special case of a module is given by $\mathcal{M} = \mathcal{A}$ (equality as sets, not as algebraic structures) with the operation in (3) being the ordinary multiplication in \mathcal{A} . This module is called the *regular module*.

Spectrum, Frequency Response, Fourier Transform: For every given \mathcal{A} and \mathcal{M} , there is an associated notion of spectrum, frequency response, and Fourier transform (if they exist). See [1] for details.

Signal Model: In applications, signals do not arise as elements of modules, but, in the discrete case considered here, as infinite or finite sequences of numbers, e.g., $\mathbf{s} = (\dots, s_{-1}, s_0, s_1, \dots) \in \mathbb{C}^{\mathbb{Z}}$ or $\mathbf{s} = (s_0, \dots, s_{n-1}) \in \mathbb{C}^n$. The purpose of the signal model, introduced next, is to assign a filter algebra \mathcal{A} and an \mathcal{A} -module \mathcal{M} to such sequences. This way, filtering is automatically defined (the operation of \mathcal{A} on \mathcal{M}), and we get access to the associated notion of spectrum and Fourier transform. In the definition, we assume complex signals, but other base fields can be chosen.

Definition 1 (Signal Model): Let $V \subseteq \mathbb{C}^I, I \subset \mathbb{Z}$, be a vector space. A *signal model* for V is a triple $(\mathcal{A}, \mathcal{M}, \Phi)$, where \mathcal{A} is an algebra, \mathcal{M} is an \mathcal{A} -module, and Φ is a bijective (one-to-one and onto) linear mapping

$$\Phi : V \rightarrow \mathcal{M}, \quad \mathbf{s} \mapsto s \in \mathcal{M}.$$

Example: Discrete Infinite Time: The abstract definition of the signal model is best illustrated by an example. Namely, the

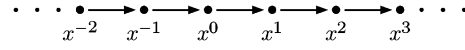


Fig. 1. Visualization of the infinite discrete time model (4) ($x = z^{-1}$).

signal model commonly adopted for infinite length discrete time SP is given by (we set $x = z^{-1}$)

$$\begin{aligned} \mathcal{A} &= \left\{ \sum_{n \in \mathbb{Z}} h_n x^n \mid \mathbf{h} = (\dots, h_{-1}, h_0, h_1, \dots) \in \ell^1(\mathbb{Z}) \right\} \\ \mathcal{M} &= \left\{ \sum_{n \in \mathbb{Z}} s_n x^n \mid \mathbf{s} = (\dots, s_{-1}, s_0, s_1, \dots) \in \ell^2(\mathbb{Z}) \right\} \\ \Phi &= \ell^2(\mathbb{Z}) \rightarrow \mathcal{M}, \quad \mathbf{s} \mapsto s = \sum_{n \in \mathbb{Z}} s_n x^n. \end{aligned} \quad (4)$$

The symbols $\ell^1(\mathbb{Z})$ and $\ell^2(\mathbb{Z})$ represent the set of infinite-length absolute summable and square summable (finite energy) sequences, respectively. As defined, $(\mathcal{A}, \mathcal{M}, \Phi)$ is a signal model for $V = \ell^2(\mathbb{Z})$ and Φ is just the ordinary z -transform. Note that in ASP $\Phi(\mathbf{s}) = s$ in (4) is primarily viewed as a formal series and not as a function. The idea is that Φ provides a basis for the coordinates \mathbf{s} and gives convolution its desired form.

Shift and Shift-Invariance: In the algebraic theory, the *shift* (or shifts) is the chosen generator (or generators) of the filter algebra. This means that every filter can be expressed as a series or polynomial in the shift (or shifts). A signal model $(\mathcal{A}, \mathcal{M}, \Phi)$ has the *shift-invariance property* if and only if \mathcal{A} is commutative. For example, the infinite discrete time model in (4) is shift-invariant, since the multiplication of Laurent series in \mathcal{A} is commutative.

Visualization: Every (discrete) signal model implicitly fixes a basis of \mathcal{M} via Φ , such as $b = (\dots, x^{-1}, x^0, x^1, \dots)$ for the time model (4). The operation of the shift on this basis can be represented by a graph, which is called the *visualization* of the model (see [1] for a rigorous definition). The visualization of (4) is shown in Fig. 1. Intuitively, it is the structure imposed by the model on the signal values \mathbf{s} , which are associated with the nodes of the graph.

B. Finite Shift-Invariant Signal Models

We identify possible signal models for finite-length 1-D sequences $\mathbf{s} = (s_0, \dots, s_{n-1}) \in V = \mathbb{C}^n$. In this case, $\dim(\mathcal{M}), \dim(\mathcal{A}) < \infty$. If we require shift-invariance (i.e., \mathcal{A} is commutative) and assume one shift, then \mathcal{A} must be a *polynomial algebra* in one variable:

$$\mathcal{A} = \mathbb{C}[x]/p(x) = \{q(x) \in \mathbb{C}[x] \mid \deg(q) < \deg(p)\}. \quad (5)$$

Here, p is an arbitrary but fixed polynomial, and addition and multiplication in \mathcal{A} is defined modulo p . The shift in \mathcal{A} is x .

In the following, we discuss signal models built from polynomial algebras and show the finite time model as an example. See [1] for more details. A good reference on polynomial algebras is [14].

Signal Model: We focus on a specific class of finite shift-invariant 1-D signal models, namely, \mathcal{A} chosen as in (5), $\mathcal{M} = \mathcal{A}$ the regular \mathcal{A} -module, and we assume that p is separable, i.e.,

has pairwise distinct zeros $\alpha = (\alpha_0, \dots, \alpha_{n-1})$. If we choose a basis $b = (p_0, \dots, p_{n-1})$ of polynomials in \mathcal{M} , then

$$\begin{aligned} \mathcal{A} &= \mathcal{M} = \mathbb{C}[x]/p(x), \\ \Phi &= \mathbb{C}^n \rightarrow \mathcal{M}, \quad \mathbf{s} \mapsto s = \sum_{0 \leq \ell < n} s_\ell p_\ell \end{aligned} \quad (6)$$

defines a signal model for $V = \mathbb{C}^n$. Filtering in this model is the multiplication $h(x)s(x) \bmod p(x)$ for $h \in \mathcal{A}$, $s \in \mathcal{M}$.

Signal Extension: Finite signals often arise because only a finite number of signal samples are available. How a finite signal continues beyond its domain is its signal extension.

Definition 2 (Signal Extension): Let $I \subset \mathbb{Z}$ and let $\mathbf{s} = (s_\ell)_{\ell \in I} \in \mathbb{C}^I$. A (linear) signal extension of \mathbf{s} is a series of linear combinations

$$s_k = \sum_{\ell \in I} \beta_{k,\ell} s_\ell, \quad \text{for } k \notin I. \quad (7)$$

If each summand contains at most one term, the signal extension is called *monomial*.

If we assume that the basis polynomials p_ℓ from b in (6) are part of an infinite sequence $(p_k)_{k \in \mathbb{Z}}$, then (6) implicitly defines a signal extension for $\mathbf{s} \in \mathbb{C}^n$. It is given by reducing p_k modulo p and expressing the result in b : $p_k \equiv \sum_{0 \leq \ell < n} \beta_{k,\ell} p_\ell$. Replacing p by s yields the signal extension in (7).

Spectrum, Fourier Transform, and Frequency Response: For the signal model (6), the spectral decomposition of \mathcal{M} , i.e., the Fourier transform, is given by the Chinese remainder theorem as

$$\begin{aligned} \Delta : \mathbb{C}[x]/p(x) &\rightarrow \mathbb{C}[x]/(x - \alpha_0) \oplus \dots \oplus \mathbb{C}[x]/(x - \alpha_{n-1}) \\ s = s(x) &\mapsto (s(\alpha_0), \dots, s(\alpha_{n-1})). \end{aligned} \quad (8)$$

$\Delta(s)$ is the spectrum of s . Further, Δ is linear¹; hence, if we choose b [which is fixed by Φ in (6)] as basis of \mathcal{M} and $x^0 = 1$ as basis in each spectral component $\mathbb{C}[x]/(x - \alpha_k)$, Δ is represented by the *polynomial transform matrix*

$$\mathcal{P}_{b,\alpha} = [p_\ell(\alpha_k)]_{0 \leq k, \ell < n}. \quad (9)$$

An arbitrary choice of bases $a_k x^0$, $a_k \in \mathbb{C}$, in the spectral components yields a *scaled polynomial transform*

$$\text{diag}(1/a_0, \dots, 1/a_{n-1}) \cdot \mathcal{P}_{b,\alpha}. \quad (10)$$

Any (scaled or not) polynomial transform is a Fourier transform for the signal model (6) and denoted with \mathcal{F} .

For a filter $h \in \mathcal{A}$, $(h(\alpha_0), \dots, h(\alpha_{n-1}))$ is the frequency response of h . Filtering $h \cdot s \pmod{p}$ is equivalent to the pointwise multiplication $(h(\alpha_0)s(\alpha_0), \dots, h(\alpha_{n-1})s(\alpha_{n-1}))$ in the spectral domain.

Filtering and Diagonalization Properties: For every filter $h \in \mathcal{A}$, filtering is a linear mapping on \mathcal{M} ; thus, with respect to the basis $b = (p_0, \dots, p_{n-1})$ of \mathcal{M} fixed by the model (6), h is represented by an $n \times n$ matrix M_h . The mapping

$$\phi : \mathcal{A} \rightarrow \mathbb{C}^{n \times n}, \quad h \mapsto \phi(h) = M_h$$

is called the *representation* of \mathcal{A} afforded by \mathcal{M} with basis b . In particular, $\phi(x)$ is called the *shift matrix*. Filtering $h \cdot s$ becomes in coordinates the matrix-vector product $\phi(h)\mathbf{s}$.

¹More precisely an \mathcal{A} -module homomorphism.



Fig. 2. Visualization of the finite discrete time model (12).

The matrices $\phi(h)$ are precisely those diagonalized by any Fourier transform \mathcal{F} for the model. Specifically,

$$\mathcal{F}\phi(h)\mathcal{F}^{-1} = \text{diag}(h(\alpha_0), \dots, h(\alpha_{n-1})). \quad (11)$$

Visualization: The graph with adjacency matrix $\phi(x)$ (the shift matrix) is the visualization of the model (6).

Example: Discrete Finite Time: As an example we consider the commonly adopted signal model for discrete finite time, given by

$$\begin{aligned} \mathcal{A} &= \mathcal{M} = \mathbb{C}[x]/(x^n - 1), \\ \Phi &= \mathbb{C}^n \rightarrow \mathcal{M}, \quad \mathbf{s} \mapsto s = \sum_{0 \leq \ell < n} s_\ell x^\ell. \end{aligned} \quad (12)$$

We call Φ the finite z -transform. Note that the chosen basis (via Φ) is $b = (x^0, \dots, x^{n-1})$. Filtering in this model is polynomial multiplication $h(x)s(x)$ modulo $(x^n - 1)$, which is equivalent to the circular convolution of \mathbf{h} and \mathbf{s} . The signal extension is obtained by reducing $x^k \equiv x^{k \bmod n} \pmod{(x^n - 1)}$ and is hence periodic and thus monomial.

The (polynomial) Fourier transform for the model (12) is readily computed via (9) as the DFT

$$\mathcal{F} = \mathcal{P}_{b,\alpha} = \text{DFT}_n = [\omega_n^{k\ell}]_{0 \leq k, \ell < n}, \quad \omega_n = e^{-2\pi\sqrt{-1}/n}.$$

For a filter $h \in \mathcal{A}$ the matrix $\phi(h)$ is a circulant matrix, which confirms the well-known property

$$\text{DFT}_n \phi(h) \text{DFT}_n^{-1} = \text{diag}(h(\omega_n^0), \dots, h(\omega_n^{n-1})).$$

The shift matrix $\phi(x)$ is the circular shift:

$$\phi(x) = \begin{bmatrix} & & & & 1 \\ & & & & \\ & & & & \\ & & & & \\ 1 & & & & \end{bmatrix}.$$

Thus, the visualization of the discrete finite time model is given by the directed circle in Fig. 2 that also captures the periodic signal extension. In words, applying a DFT to a signal $\mathbf{s} \in \mathbb{C}^n$ associates the values s_ℓ with the nodes of this graph, which is equivalent to imposing a periodic signal extension.

C. Derivation of Signal Models

In [1], we presented a procedure to derive infinite and finite signal models from an abstract definition of the shift operation. We used this procedure to derive the infinite and finite time models (4) and (12) from the standard time shift

$$\text{time shift : } q \diamond t_n = t_{n+1} \quad n \in \mathbb{Z}. \quad (13)$$

displayed in Fig. 3 (top). Here the t_n denote abstract time marks, q is the shift operator, and \diamond is the shift operation.

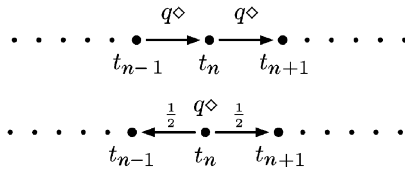


Fig. 3. Time shift (top) and space shift (bottom).

The procedure consists of three steps. First, the shift is defined in the abstract form shown in (13) and a k -fold shift q_k is introduced through $q_k \diamond t_n = t_{n+k}$. This implies that $q_k = q^k$. Second, the shift operation is extended to linear combinations $\sum s_n t_n$ of the time marks and to linear combinations of k -fold shifts $q_k: \sum h_k q_k$. Third, the model is realized by setting $q = x$, replacing \diamond with ordinary multiplication, and solving

$$t_{n+1} = x \cdot t_n \tag{14}$$

for t_n . Normalizing $t_0 = 1$ yields $t_n = x^n$ as unique solution. In the infinite case, convergence requirements lead to the model in (4). In the finite case, as was shown in [1], a boundary condition is needed to ensure that \mathcal{M} becomes a module. This boundary condition determines the entire signal extension, and requiring a monomial signal extension (the simplest possible; see Definition 2) leads to $\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/(x^n - a)$, $a \in \mathbb{C}$. For $a = 1$ this yields the finite time model (12).

In the following sections, we derive signal models for discrete infinite and finite space. These models are built using the same procedure but starting from a different definition of the shift.

III. INFINITE 1-D SPACE MODELS

Standard SP considers time-invariant systems, which implies the standard definition of the shift in (13). In this section and the next, we will use ASP to derive an SP framework for space SP as we refer to it. It is built from a different, symmetric definition of the shift. We have two motivations for this definition. The first is our goal to define the shift for signals for which there is no intrinsic sense of direction. These signals contrast with time signals, for which past, present, and future are inherent from the direction of time. The second reason is, as we will show, that our space shift definition leads to signal models that have the 16 DTTs as Fourier transforms. Thus, within ASP, time and space SP, the DFT and the DTTs become instantiations of one general framework. There will be many other benefits of this theoretical exercise as discussed later.

A. Constructing the Signal Model

We follow the same steps as in the time model derivation in [1].

Definition of the Shift: We consider discrete complex signals indexed by $\mathbb{Z}: s \in \mathbb{C}^{\mathbb{Z}}$; i.e., we consider the vector space $V = \mathbb{C}^{\mathbb{Z}}$. We define now *space* marks t_n and an appropriate *space* shift operator q and its operation \diamond on the space marks. As mentioned above, q should operate symmetrically. We adopt the definition

$$\text{space shift : } q \diamond t_n = (t_{n+1} + t_{n-1})/2, \quad n \in \mathbb{Z} \tag{15}$$

visualized in Fig. 3 (bottom).

We proceed by extending the operator domain from q to k -fold shift operators q_k . A natural definition of the k -fold space shift is

$$q_k \diamond t_n = (t_{n+k} + t_{n-k})/2 \tag{16}$$

since t_{n+k} and t_{n-k} are those space marks at distance k from t_n .

Here, we have the first interesting difference with respect to the time model derivation, since clearly $q_k \neq q^k$. Furthermore, (16) implies $q_k = q_{-k}$; hence, it is sufficient to consider only shift operators q_k with $k \geq 0$. Thus, the natural representation of a filter will be $\sum_{k \geq 0} h_k q_k$. The following lemma shows that the q_k are given by the Chebyshev polynomials of the first kind T_k in the variable q . The Chebyshev polynomials will play a central role in the definition of space models. For this reason, we provide the necessary background on four types of Chebyshev polynomials T, U, V , and W in Appendix I, which we encourage the reader to briefly review at this point.

Lemma 3: The k -fold space shift operator is given by $q_k = T_k(q)$.

Proof: Induction on k . By definition $q_0 = 1$, and $q_1 = q = T_1(q)$. Also by definition, $q_{k+1} \diamond t_n = (t_{n+k+1} + t_{n-k-1})/2 = (t_{n+k+1} + t_{n+k-1} + t_{n-k+1} + t_{n-k-1})/2 - (t_{n+k-1} + t_{n-k+1})/2 = 2q \diamond (t_{n+k} + t_{n-k})/2 - (t_{n+k-1} + t_{n-k+1})/2 = (2qq_k - q_{k-1}) \diamond t_n$, for $n \in \mathbb{Z}$. From the induction hypothesis, $q_k = T_k(q)$, $q_{k-1} = T_{k-1}(q)$, and thus, using the recurrence of the Chebyshev polynomials [(43) in Appendix I], $q_{k+1} = T_{k+1}(q)$, as desired. ■

Linear Extension: To construct a linear signal model, we extend by linearity the operation of q to the entire set $\mathcal{M} = \{s = \sum_{n \in \mathbb{Z}} s_n t_n\}$, namely as $q \diamond s = \sum_{n \in \mathbb{Z}} s_n (q \diamond t_n)$, which can be evaluated. Similarly, we linearly extend the operator domain to $\mathcal{A} = \{h = \sum_{k \geq 0} h_k q_k\} = \{h = \sum_{k \geq 0} h_k T_k(q)\}$ using Lemma 3.

Realization: We determine a “realization” of the model introduced in the previous section. We set in (15) $q = x$, $\diamond = \cdot$, and determine polynomials C_n that replace the space marks t_n in (15), i.e., that satisfy

$$x \cdot C_n = (C_{n+1} + C_{n-1})/2. \tag{17}$$

Since (17) is equivalent to (43) (in Appendix I), the solution is given by a sequence of Chebyshev polynomials.

We immediately notice differences with respect to the corresponding derivation in the time case. These differences are intrinsic to the space model.

- Equation (17) is a three-term recurrence for the space marks, whereas (14) is a two-term recurrence for the time marks.
- Only the C_n , $n \geq 0$, are linearly independent; the C_n , $n < 0$, are polynomials in x and can thus be expressed as linear combinations of $\{C_n | n \geq 0\}$. In other words, the realization of the space model introduces a starting point in space, given by $C_0 = 1$. Fixing C_1 determines the left boundary condition and the left signal extension.
- As a consequence, even after normalizing $C_0 = 1$, the sequence C_n of Chebyshev polynomials is not uniquely determined. The degree of freedom is given by the choice of C_1 as a polynomial of degree 1.

TABLE I
REALIZATION OF THE ABSTRACT SPACE MODEL

concept	abstract	realized
shift operator	q	$T_1(x) = x$
shift operation	\diamond	\cdot
space mark	t_n	C_n
k -fold shift	$q_k = T_k(q)$	$T_k(x)$
space shift	$q \diamond t_n = \frac{1}{2}(t_{n+1} + t_{n-1})$	$x \cdot C_n = \frac{1}{2}(C_{n+1} + C_{n-1})$
signal	$\sum s_n t_n$	$\sum s_n C_n(x)$
filter	$\sum h_k T_k(q)$	$\sum h_k T_k(x)$

- Again, we note that in the time model, a k -fold shift operator is given by x^k :

$$x^k \cdot x^n = x^{n+k}$$

in contrast to the space model, where, by Lemma 3, the k -fold shift operator is given by $T_k(x)$, independent of C [see Lemma 14(iv) in Appendix I]:

$$T_k \cdot C_n = (C_{n+k} + C_{n-k})/2. \tag{18}$$

As a result of this discussion, we obtain the spaces $\mathcal{A} = \{h = \sum_{k \geq 0} h_k T_k\}$ and $\mathcal{M} = \{s = \sum_{n \geq 0} s_n C_n\}$, i.e., the signal model that we obtain later will be only for right-sided sequences.

Table I shows the correspondence between abstract and realized concepts.

To ensure convergence, we would like to require as before $\mathbf{h} \in \ell^1(\mathbb{N})$ and $\mathbf{s} \in \ell^2(\mathbb{N})$. However, to prove convergence, we have first to choose proper boundary conditions, i.e., we have to choose the proper Chebyshev polynomials C . We analyze the boundary conditions in the next paragraph. This discussion has no counterpart in the derivation of the infinite time model in [1].

Left Boundary Condition and Left Signal Extension: The degree of freedom for choosing a Chebyshev sequence C , normalized by $C_0 = 1$, is given by the choice of C_1 , or, equivalently, by the choice of C_{-1} , since the entire sequence is then obtained by applying the Chebyshev recursion (43) in both directions [see Lemma 14(i) in Appendix I]. Fixing either C_1 or C_{-1} is equivalent to choosing a *left boundary condition* for the signal $\mathbf{s} = (s_0, s_1, \dots)$. For example, setting $C_1 = x$ implies $C_{-1} = x$, and thus $C_{-1} = C_1$, which imposes on the signal \mathbf{s} the left boundary condition $s_{-1} = s_1$. Using Table VII, the corresponding sequence is $C = T$.

To determine the left boundary condition in the general case, we set $C_0 = 1$ and $C_1 = ax + b$, $a \neq 0$ (to satisfy $\deg(C_1) = 1$). Then, by applying (43) backwards, we get

$$C_{-1} = 2x - (ax + b) = \frac{2-a}{a}C_1 - \frac{2b}{a}C_0. \tag{19}$$

Since C_{-1} is of degree at most 1, every polynomial C_{-n} , $n > 0$, obtained by the recursion (43), is of degree at most n , and thus a linear combination of the polynomials C_0, \dots, C_n ,

$$C_{-n} = \sum_{0 \leq i \leq n} \beta_i \cdot C_i. \quad n > 0 \tag{20}$$

This equation defines the *left signal extension* associated with the sequence C . On the other hand, by comparing the degrees

of freedom, it is obvious that not every signal extension can be obtained by choosing a suitable boundary condition. Thus

$$(C \Leftrightarrow \text{left boundary condition}) \Rightarrow \text{left signal extension.}$$

For a generic left boundary condition, the left signal extension (20) has no simple structure; in particular, it is not monomial (see Definition 2). We determine now those left boundary conditions that yield a *monomial* left signal extension in (20). The answer is provided in the following lemma.

Lemma 4 (Monomial Left Signal Extension): Let $C = (C_n | n \in \mathbb{Z})$ be a sequence of Chebyshev polynomials with $C_0 = 1$ and $\deg(C_1) = 1$. Then the left signal extension associated with C is monomial, i.e., every C_k , $k < 0$, is a multiple of a C_n , $n \geq 0$, if and only if $C \in \{T, U, V, W\}$ (see Appendix I), i.e., $C_{-1} \in \{C_1, 0, C_0, -C_0\}$, which implies the corresponding left boundary conditions $s_{-1} \in \{s_1, 0, s_0, -s_0\}$.

Proof: If $C \in \{T, U, V, W\}$, then the assertion holds as shown in the ‘‘symmetry’’ column of Table VII. It remains to show the converse. We start with the generic left boundary condition in (19). Because the signal extension associated with C is monomial, one of the two summands in (19) has to vanish.

Case 1: C_{-1} is a multiple of C_0 , i.e., constant. It follows $a = 2$, $C_1 = 2x + b$, $C_{-1} = -b$, $C_{-2} = -2bx - 1$. Now, either C_{-2} is constant, i.e., $b = 0$, which implies $C = U$, or C_{-2} is a multiple of C_1 , which implies $b = \pm 1$, or $C \in \{V, W\}$.

Case 2: C_{-1} is a multiple of C_1 . It follows $b = 0$, $C_1 = ax$, $C_2 = 2ax^2 - 1$, $C_{-1} = (2 - a)x$, $a \neq 2$, and $C_{-2} = 2(2 - a)x^2 - 1$. Since C_{-2} has to be a multiple of C_2 , we get $a = 1$ and thus $C = T$. This completes the proof. ■

The four boundary conditions in Lemma 4 are the discrete versions of the so-called Dirichlet boundary condition (‘‘zero value’’) and von Neumann boundary condition (‘‘zero slope’’) [15], [16]. In each case, the symmetry point is either a ‘‘whole’’ sample point, or a ‘‘half’’ sample point, i.e., is located between two sample points. In the literature, these four signal extensions are sometimes called: whole point symmetry (WS), whole point antisymmetry (WA), half point symmetry (HS), and half point antisymmetry (HA) [17].

For these four choices of boundary conditions, filtering, i.e., the multiplication $\sum_{n \geq 0} h_n T_n \cdot \sum_{n \geq 0} s_n C_n$ converges provided $\mathbf{h} \in \ell^1(\mathbb{N})$, $\mathbf{s} \in \ell^2(\mathbb{N})$ (see [6] for more details).

Resulting Infinite Space Models: We define four infinite space models $(\mathcal{A}, \mathcal{M}, \Phi)$ for $V = \ell^2(\mathbb{N})$. Namely, for $C \in \{T, U, V, W\}$

$$\begin{aligned} \mathcal{A} &= \left\{ h = \sum_{k \geq 0} h_k T_k(x) \mid \mathbf{h} \in \ell^1(\mathbb{N}) \right\}, \\ \mathcal{M} &= \left\{ s = \sum_{n \geq 0} s_n C_n(x) \mid \mathbf{s} \in \ell^2(\mathbb{N}) \right\}, \\ \Phi : \ell^2(\mathbb{N}) &\rightarrow \mathcal{M}, \mathbf{s} \mapsto \sum_{n \geq 0} s_n C_n(x). \end{aligned} \tag{21}$$

We call Φ the C -transform but will replace C by either T , U , V , or W , when appropriate, and accordingly refer to the T -, U -, V -, or W -transform.

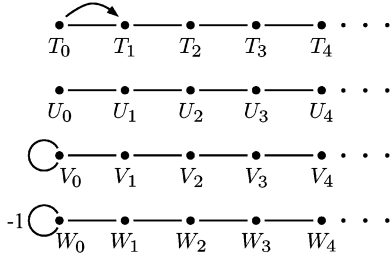


Fig. 4. Visualization of the four infinite space models for $C \in \{T, U, V, W\}$. The common edge scaling factor $1/2$ is omitted.

B. Properties

Each of these models has its associated notion of filtering, spectrum, frequency response, and Fourier transform as explained in [1]. We omit the details here since our focus are the finite space models that we will show to underly the DTTs.

The visualizations of the models are shown in Fig. 4 with a common scaling factors of $1/2$ omitted. The graphs are undirected, since they are space models. Namely, the space shift (Fig. 3 bottom) yields between each two space marks an edge in both directions. The behavior at the left edge is determined by the left boundary condition. Namely, $x C_0 = \frac{1}{2}(C_{-1} + C_1)$ produces a directed edge to the (nonexistent) C_{-1} . In the first case, $C = T, T_{-1} = T_1$, and hence this edge is rerouted to T_1 . In the second case, $C = U, U_{-1} = 0$; hence, the edge vanishes.

IV. FINITE 1-D SPACE MODELS AND DTTs

In this section, we derive finite versions of the space models in (21). As in the finite time model (12), these space models will have polynomial algebras as filter and signal spaces. This is not surprising as ASP explains that only those choices support shift-invariance (Section II-A). We derive the finite space models in the same way as we derived the finite time model in [1], namely by requiring a monomial signal extension. However, in contrast with the time case, this signal extension will not be periodic but symmetric or antisymmetric with 16 choices. This is due to the different basis required after realizing the shift operation: x^ℓ supports the time shift, C_ℓ supports the space shift.

By applying the general theory from Section II-B, we will see that the Fourier transforms for the finite space models are precisely the 16 DTTs. There are various benefits to knowing these models. First, as application of the general theory in Section II-B, we obtain the appropriate notions of “z-transform,” filtering or convolution, convolution theorems, spectrum and frequency response associated with the DTTs and can derive and explain many of their properties. Second, we establish that the DTTs are, in a rigorous sense, associated with the space shift, Fig. 3 (bottom), in the same way as the DFT is associated with the time shift. Third, knowing those signal models is the key to deriving and understanding the DTTs’ fast algorithms [7], [9].

A. Constructing the Signal Model

Shift, Linear Extension, Realization: We consider a finite number of space marks t_0, \dots, t_{n-1} and adopt the space shift operator q in Fig. 3 (bottom) and its realization by setting $q = x$,

and hence $t_k = C_k$ (a generic sequence of Chebyshev polynomials),² as derived in Section III-A. These definitions will need to be complemented by appropriate boundary conditions, as we discuss next.

Let $\mathbf{s} = (s_0, \dots, s_{n-1}) \in \mathbb{C}^n$ be a finite sampled signal and C a sequence of Chebyshev polynomials. A straightforward realization seems to lead to signals that are polynomials of the form $\sum_{0 \leq k < n} s_k C_k$. The set of these is the vector space $\mathbb{C}_n[x]$ of polynomials of degree less than n (with basis polynomials C_k). However, this space is not closed under multiplication by the shift operator x , and thus it is not a module, which means filtering is not well-defined. In particular, the problem is that

$$x \cdot C_{n-1} = (C_{n-2} + C_n)/2 \notin \mathbb{C}_n[x] \tag{22}$$

since $C_n \notin \mathbb{C}_n[x]$. Note that, in contrast to the time case [1], the left boundary does not impose any problems, since

$$x \cdot C_0 = (C_{-1} + C_1)/2 \in \mathbb{C}_n[x].$$

Namely, the choice of C already implies a left boundary condition via (19). So the remaining task is to determine the proper right boundary conditions.

Right Boundary Condition and Signal Extension: To solve the problem in (22), we introduce an equation

$$C_n = r = \sum_{0 \leq k < n} \beta_k C_k, \quad \text{or} \quad C_n - r = 0. \tag{23}$$

This imposes the same equation on the corresponding signal samples s_k associated with C_k , namely

$$s_n = \sum_{0 \leq k < n} \beta_k s_k$$

which is the right boundary condition. As a consequence of (23), using the k -fold space shift operator T_k (see Lemma 3), we get the series of equations for $k \geq 0$

$$0 = 2T_k \cdot 0 \equiv 2T_k(C_n - r) = C_{n+k} + C_{n-k} - 2T_k r$$

which determine the entire right signal extension. It is obtained by reducing $C_{n+k} \equiv 2T_k r - C_{n-k}$ modulo $(C_n - r)$.

Algebraically, the right boundary condition replaces the vector space $\mathbb{C}_n[x]$ (with basis $b = (C_0, \dots, C_{n-1})$) by $\mathcal{M} = \mathbb{C}[x]/(C_n - r)$ (also with basis b), viewed as a regular module, i.e., the algebra is $\mathcal{A} = \mathcal{M}$. The natural basis in \mathcal{A} is given by (T_0, \dots, T_{n-1}) , regardless of the choice of C .

For a general choice of left boundary condition (given by the choice of C) and right boundary condition (given by the choice of r), the corresponding signal extension has a complicated structure. As before, we identify those boundary conditions that lead to a monomial signal extension. Lemma 4 gives already the proper left boundary conditions and shows that they are obtained by choosing $C \in \{T, U, V, W\}$. For the right boundary conditions, there are again four choices, which yields a total number of 16 possibilities—corresponding to the 16 types of DTTs as we will see below.

²We note that another realization is possible by setting q not equal to x . However, the derived space models have two-dimensional spectral components, which is undesirable. See [6] for details.

TABLE II
THE 16 POLYNOMIALS p ASSOCIATED WITH THE 16 FINITE SPACE MODELS.
 C_n HAS TO BE REPLACED BY T_n, U_n, V_n, W_n TO OBTAIN ROWS
1,2,3,4, RESPECTIVELY

C_n	$C_n - C_{n-2}$	C_n	$C_n - C_{n-1}$	$C_n + C_{n-1}$
T_n	$2(x^2 - 1)U_{n-2}$	T_n	$(x - 1)W_{n-1}$	$(x + 1)V_{n-1}$
U_n	$2T_n$	U_n	V_n	W_n
V_n	$2(x - 1)W_{n-1}$	V_n	$2(x - 1)U_{n-1}$	$2T_n$
W_n	$2(x + 1)V_{n-1}$	W_n	$2T_n$	$2(x + 1)U_{n-1}$

Lemma 5 (Monomial Right Signal Extension): For a monomial left signal extension, let $C \in \{T, U, V, W\}$. The only four right boundary conditions that yield a monomial signal extension for $\mathcal{M} = \mathbb{C}[x]/p(x)$ are $C_n = C_{n-2}, C_n = 0$, and $C_n = \pm C_{n-1}$, which implies $p \in \{C_n - C_{n-2}, C_n, C_n \pm C_{n-1}\}$. These 16 p 's are shown in Table II.

Proof: Necessarily, the boundary condition has the form $C_n = aC_k, 0 \neq k < n$. By multiplying by x on both sides, we obtain $C_{n+1} = a(C_{k+1} + C_{k-1}) - C_{n-1}$. We determine under which conditions the three summands on the right reduce to at most one summand.

Case 1: $k \neq n - 1$: Then either $a = 0$, or $k = n - 2$ and $a = 1$.

Case 2: $k = n - 1$: Then $aC_{k+1} = aC_n = a^2C_{n-1}$ and thus $a = \pm 1$.

It remains to show that these four boundary conditions yield a monomial signal extension, which is done by induction. We omit the details.

The identities in Table II are obtained using Table VII in Appendix I and well-known trigonometric identities. ■

It is interesting to note that the right boundary conditions in Lemma 5 are the reflections of the left boundary conditions in Lemma 4.

Resulting Finite Space Models: We define 16 finite space models $(\mathcal{A}, \mathcal{M}, \Phi)$ for $V = \mathbb{C}^n$. Namely, for $C \in \{T, U, V, W\}$ and $p \in \{C_n - C_{n-2}, C_n, C_n \pm C_{n-1}\}$

$$\begin{aligned} \mathcal{A} &= \mathbb{C}[x]/p(x) = \left\{ h = \sum_{0 \leq k < n} h_k T_k(x) \mid h_k \in \mathbb{C} \right\} \\ \mathcal{M} &= \mathbb{C}[x]/p(x) = \left\{ s = \sum_{0 \leq k < n} s_k C_k(x) \mid s_k \in \mathbb{C} \right\} \\ \Phi : \mathbb{C}^n &\rightarrow \mathcal{M}, \mathbf{s} \mapsto \sum_{0 \leq k < n} s_k C_k(x). \end{aligned} \tag{24}$$

We call each Φ a finite C -transform, and replace C with T, U, V , or W if specified. Note that $\mathcal{A} = \mathcal{M}$ but the natural basis in \mathcal{A} always consists of the k -fold space shifts T_k , independently of C .

Example: We choose the left boundary condition $s_{-1} = s_0$, i.e., $C_{-1} = C_0$, which is afforded by the base polynomials $C = V$. As right boundary condition, we choose $s_n = s_{n-1}$, i.e., $C_n = C_{n-1}$, which implies

$$p = C_n - C_{n-1} = V_n - V_{n-1} = 2(x - 1)U_{n-1}$$

TABLE III
EIGHT TYPES OF DCTS AND DSTS (UNSCALED) OF SIZE n . THE ENTRY AT
ROW k AND COLUMN ℓ IS GIVEN FOR $0 \leq k, \ell < n$

type	DCT	DST
1	$\cos k\ell \frac{\pi}{n-1}$	$\sin(k+1)(\ell+1) \frac{\pi}{n+1}$
2	$\cos k(\ell + \frac{1}{2}) \frac{\pi}{n}$	$\sin(k+1)(\ell + \frac{1}{2}) \frac{\pi}{n}$
3	$\cos(k + \frac{1}{2})\ell \frac{\pi}{n}$	$\sin(k + \frac{1}{2})(\ell + 1) \frac{\pi}{n}$
4	$\cos(k + \frac{1}{2})(\ell + \frac{1}{2}) \frac{\pi}{n}$	$\sin(k + \frac{1}{2})(\ell + \frac{1}{2}) \frac{\pi}{n}$
5	$\cos k\ell \frac{\pi}{n-\frac{1}{2}}$	$\sin(k+1)(\ell+1) \frac{\pi}{n+\frac{1}{2}}$
6	$\cos k(\ell + \frac{1}{2}) \frac{\pi}{n-\frac{1}{2}}$	$\sin(k+1)(\ell + \frac{1}{2}) \frac{\pi}{n+\frac{1}{2}}$
7	$\cos(k + \frac{1}{2})\ell \frac{\pi}{n-\frac{1}{2}}$	$\sin(k + \frac{1}{2})(\ell + 1) \frac{\pi}{n+\frac{1}{2}}$
8	$\cos(k + \frac{1}{2})(\ell + \frac{1}{2}) \frac{\pi}{n-\frac{1}{2}}$	$\sin(k + \frac{1}{2})(\ell + \frac{1}{2}) \frac{\pi}{n-\frac{1}{2}}$

using Table II. We obtain the associated signal model (the 2 in p can be dropped)

$$\begin{aligned} \mathcal{A} &= \mathcal{M} = \mathbb{C}[x]/(x - 1)U_{n-1}(x) \\ \Phi : \mathbf{s} &\mapsto s = \sum_{0 \leq \ell < n} s_\ell V_\ell. \end{aligned} \tag{25}$$

We will see later that the DCT, type 2, is a Fourier transform for this model.

Next, we apply the general theory from Section II-B to all 16 finite space models.

B. Spectrum and Fourier Transform: DTTs

We show that the 16 DTTs are Fourier transforms for the 16 finite space models (24). In doing so, we settle the question why there are 16 DTTs to begin with, as the original derivation of the full set of all 16 [4] does not provide an explanation.

The first and most important DTT is the DCT, type 2, introduced in [2] and used in the JPEG image compression standard. Table III gives the definitions of the *nonorthogonal* versions of the 16 DTTs. We note that the DTTs of type 1, 4, 5, and 8 are symmetric, and that the DTTs of type 2 and 3, 6 and 7, respectively, are transposes of each other. We use Arabic instead of Roman numbers to denote the type following [16].

To compute the Fourier transform (8) of the finite space models (24) and its matrix form \mathcal{F} in (9) or (10), we have to determine the zeros of the 16 polynomials in Table II, which can be done using Table VII in Appendix I. Instead of giving the details for all 16 cases, we consider the signal model (25) as a representative example and then state the result for all 16 DTTs. Note that the discussion is an application of the general theory in Section II-B.

Example: DCT, Type 2: The zeros of $(x - 1)U_{n-1}(x)$ in (25) are given by $\alpha_k = \cos(k\pi/n), 0 \leq k < n$ (from Table VII in Appendix I). Hence, the Fourier transform for \mathcal{M} is given by

$$\begin{aligned} \Delta : \mathbb{C}[x]/(x-1)U_{n-1}(x) &\rightarrow \bigoplus_{0 \leq k < n} \mathbb{C}[x]/(x-\alpha_k) \\ s &= s(x) \mapsto (s(\alpha_0), \dots, s(\alpha_{n-1})). \end{aligned} \tag{26}$$

$\Delta(s)$ is the spectrum of the signal s and $(h(\alpha_0), \dots, h(\alpha_{n-1}))$ is the frequency response of the filter $h \in \mathcal{A}$.

TABLE IV
 OVERVIEW OF THE 16 DTTs AND THEIR ASSOCIATED SIGNAL MODELS. THE LEFT BOUNDARY CONDITION (ROWS) DETERMINES A SCALING FUNCTION $f(\cos \theta = x)$ AND THE BASIS $C \in \{T, U, V, W\}$ IN $\mathcal{M} = \mathbb{C}[x]/p(x)$ AND HENCE Φ . THE RIGHT BOUNDARY CONDITION (COLUMNS) THEN DETERMINES $p(x)$ (GIVEN BELOW THE DTT) AND HENCE ALSO $\mathcal{A} = \mathcal{M}$

	$s_n - s_{n-2}$	s_n	$s_n - s_{n-1}$	$s_n + s_{n-1}$	f	C
$s_{-1} = s_1$	DCT-1 $2(x^2 - 1)U_{n-2}$	DCT-3 T_n	DCT-5 $(x - 1)W_{n-1}$	DCT-7 $(x + 1)V_{n-1}$	1	T
$s_{-1} = 0$	DST-3 $2T_n$	DST-1 U_n	DST-7 V_n	DST-5 W_n	$\sin \theta$	U
$s_{-1} = s_0$	DCT-6 $2(x - 1)W_{n-1}$	DCT-8 V_n	DCT-2 $2(x - 1)U_{n-1}$	DCT-4 $2T_n$	$\cos \frac{1}{2}\theta$	V
$s_{-1} = -s_0$	DST-8 $2(x + 1)V_{n-1}$	DST-6 W_n	DST-4 $2T_n$	DST-2 $2(x + 1)U_{n-1}$	$\sin \frac{1}{2}\theta$	W

In matrix form, the unique polynomial Fourier transform (9) for the signal model has entries

$$V_\ell(\alpha_k) = \frac{1}{\cos k\pi/(2n)} \cdot \cos k(\ell + 1/2)\pi/n.$$

We can scale these to cancel the denominator and get the matrix

$$\begin{aligned} & \text{diag}_{0 \leq k < n} (\cos k\pi/(2n)) \cdot [V_\ell(\alpha_k)]_{0 \leq k, \ell < n} \\ &= [\cos k(\ell + 1/2)\pi/n]_{0 \leq k, \ell < n} \\ &= \text{DCT-}2_n. \end{aligned} \tag{27}$$

In words, the DCT- 2_n is a Fourier transform [namely a scaled polynomial transform (10)] for the signal model (25). The scaling diagonal in (27) shows the basis chosen on the right-hand side of (26), namely $1/(\cos k\pi/(2n))$ in the k th spectral component $\mathbb{C}[x]/(x - \cos k\pi/n)$, $0 \leq k < n$.

All DTTs: Similar computations for all 16 cases establishes the 16 DTTs as Fourier transforms for the 16 finite C -transforms.

Theorem 6 (DTTs and Polynomial Algebras): The 16 DTTs are the Fourier transforms for the 16 finite space models (24). The correspondence is given in Table IV as follows. Let $(\mathcal{A}, \mathcal{M}, \Phi)$ be a finite space model with $\mathcal{M} = \mathbb{C}[x]/p$ with basis $b = (C_0, \dots, C_{n-1})$. The choice of C (rows of Table IV) determines the left boundary condition and a scaling function f . The choice of right boundary condition (columns 2–5 in Table IV) then determines the polynomial p , given at the intersection of row and column. The corresponding DTT is given above p . More specifically, assume $\alpha = (\alpha_0, \dots, \alpha_{n-1})$ are the zeros of p . All α_k have the form $\cos r_k\pi$, $r_k \in [0, \pi]$ (see Table VII in the Appendix I), and α is ordered by increasing r_k . Then

$$\text{DTT}_n = \text{diag}_{0 \leq k < n} (f(\alpha_k)) \cdot \mathcal{P}_{b,\alpha} \tag{28}$$

i.e., DTT_n is a scaled polynomial transform and thus a Fourier transform for the associated signal model (see Section II-B). Equation (28) implies that the chosen basis in the spectral component $\mathbb{C}[x]/(x - \alpha_k)$ is $1/f(\alpha_k)$, $0 \leq k < n$.

The DCT, type 3, was implicitly recognized as a polynomial transform in [10]. The DCTs and DSTs of types 1–4 were recognized as (scaled) polynomial transforms in [11]. In both cases

no connection to signal processing was established. The original derivation of the DCT, type 2, in [2] mentions Chebyshev polynomials but does not make use of this fact nor connects to algebra.

Polynomial DTTs: Theorem 6 shows that each DTT is a Fourier transform for a finite space model but in general not the corresponding polynomial transform. Thus, we now associate to each DTT its polynomial transform $\mathcal{P}_{b,\alpha}$ obtained by omitting the scaling factors in (28).

Definition 7 (Polynomial DTTs): Let DTT_n be given. We call the unique polynomial transform $\mathcal{P}_{b,\alpha}$ associated with DTT_n by (28) the “polynomial DTT” and denote it with $\overline{\text{DTT}}_n$. Thus, (28) can be rewritten as

$$\text{DTT}_n = \text{diag}_{0 \leq k < n} (f(\alpha_k)) \cdot \overline{\text{DTT}}_n.$$

We have $\text{DTT} = \overline{\text{DTT}}$ if and only if DTT appears in the first row of Table IV, i.e., if $\text{DTT} \in \{\text{DCT-1}, \text{DCT-3}, \text{DCT-5}, \text{DCT-7}\}$.

The polynomial DTTs will play an important role in the derivation of fast DTT algorithms [7]. Also, in some cases the polynomial DTTs have a lower complexity than the actual DTT. This makes them a candidate for applications in which the DTT is followed by scaling (such as JPEG compression).

Remarks and Observations: For each DTT, we have three relevant versions. First, the polynomial version $\overline{\text{DTT}}$, which is the unique polynomial transform for its associated signal model (see Definition 7 above). Second, the unscaled or natural version, which has pure cosines (or sines) as entries (see Table III). Third, the orthogonal version, which arises from the other two by suitable scaling of rows and columns, i.e., by slightly adjusting the signal model (explained below in Section IV-E).

The 16 DTTs can be divided into four groups of four each with respect to the polynomial p in the associated module $\mathbb{C}[x]/p$ (see Table IV). For example, the “ T -group” comprises all DTTs of types 3 and 4, which have the same module $\mathcal{M} = \mathbb{C}[x]/T_n$. The modules within the other groups differ slightly, e.g., in the U -group that comprises the DTTs on the main diagonal in Table IV. The difference between the DTTs within the same group is the choice of basis, which is one of T, U, V, W . As a consequence, these transforms can be converted into each other using a sparse base change (explained in Section IV-F).

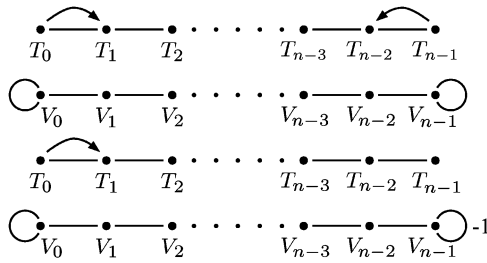


Fig. 5. Visualizations of the finite space models associated with the DCTs of type 1-4 (from top to bottom) and size n . A common edge scaling factor of $1/2$ has been omitted.

TABLE V
THE VALUES $\beta_1, \beta_2, \beta_3, \beta_4$ FROM (30) FOR THE FOUR RESPECTIVE CHOICES OF LEFT AND RIGHT BOUNDARY CONDITION

left boundary condition	β_1	β_2	right boundary condition	β_3	β_4
$s_{-1} = s_1$	0	2	$s_n = s_{n-2}$	2	0
$s_{-1} = 0$	0	1	$s_n = 0$	1	0
$s_{-1} = s_0$	1	1	$s_n = s_{n-1}$	1	1
$s_{-1} = -s_0$	-1	1	$s_n = -s_{n-1}$	1	-1

C. Visualization

The right boundary conditions for the 16 finite space models (24) are precisely the mirrored left boundary conditions that occurred already in Fig. 4. This makes it easy to obtain the visualizations for (24). For example, Fig. 5 shows the cases associated with the DCTs of type 1-4.

More formally, consider the model (25) as example. To obtain the visualization, we have to compute the shift matrix $\phi(x)$. From $xV_0 = \frac{1}{2}V_0 + \frac{1}{2}V_1, xV_i = \frac{1}{2}V_{i-1} + \frac{1}{2}V_{i+1}, 1 \leq i < n-1$, and $xV_{n-1} = \frac{1}{2}V_{n-2} + \frac{1}{2}V_n \equiv \frac{1}{2}V_{n-2} + \frac{1}{2}V_{n-1} \pmod{(V_n - V_{n-1})}$, it follows that

$$\phi(x) = \frac{1}{2} \cdot \begin{bmatrix} 1 & 1 & & & \\ 1 & 0 & \cdot & & \\ 0 & 1 & \cdot & 1 & \\ & & \cdot & 0 & 1 \\ & & & 1 & 1 \end{bmatrix}. \tag{29}$$

This is precisely the adjacency matrix of the second graph in Fig. 5 associated with the DCT of type 2. In words, applying the DCT-2 to a signal s implicitly imposes the structure of this graph on that signal.

For an arbitrary finite space model (24), $\phi(x)$ takes the form

$$\phi(x) = \frac{1}{2} \cdot \begin{bmatrix} \beta_1 & 1 & & & \\ \beta_2 & 0 & \cdot & & \\ 0 & 1 & \cdot & 1 & \\ & & \cdot & 0 & \beta_3 \\ & & & 1 & \beta_4 \end{bmatrix} \tag{30}$$

with the β_i shown in Table V.

D. Filtering and Diagonalization Properties

Consider a finite space model (24) with $\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/p(x)$ and C -basis b (fixed by Φ) and associated DTT_n . Let ϕ be the representation associated with the model.

Filtering in this model is the multiplication of $h \in \mathcal{A}$ (expressed in the T -basis) with $s \in \mathcal{M}$ (expressed in b) modulo p to yield again a signal expressed in b . In coordinates, $hs \pmod p$ is equivalent to $\phi(h)s$.

The diagonalization properties of the 16 DTTs are a special case of (11) and can be stated in a unified way. For any filter $h \in \mathcal{A}$,

$$DTT_n \cdot \phi(h) \cdot DTT_n^{-1} = \text{diag}_{0 \leq k < n}(h(\alpha_k)) \tag{31}$$

where the α_k are the zeros of $p(x)$. This unifies and explains the result from [18]. Conversely, the $\phi(h)$ are all the matrices diagonalized by DTT. The matrices $\phi(h)$ have in all cases structure: each $\phi(h)$ can be written as the sum of a Toeplitz and a Hankel matrix, up to potential scaling factors. More details are in [6].

As one example, for $\phi(x)$ in (29), we get

$$DCT\text{-}2_n \cdot \phi(x) \cdot DCT\text{-}2_n^{-1} = \text{diag}_{0 \leq k < n}(\cos(k\pi/n)).$$

More generally, the DTTs diagonalize their associated $\phi(x)$ in (30) via (31), which was also observed in [16] (where $\phi(2-2x) = 2I_n - 2\phi(x)$ was considered instead of $\phi(x)$). This also implies that the $\phi(x)$ have pairwise distinct eigenvalues.

Equation (31) also provides the convolution theorems associated with the finite space models.

E. Orthogonal DTTs

It is well known that the DTTs, as defined in Table III, are ‘‘almost orthogonal,’’ which means that after a suitable scaling of rows and columns they become orthogonal. Using ASP, i.e., the knowledge of the DTTs’ underlying signal models (24), these scaling factors can be derived as explained in [6] and omitted here due to space limitations.

Another argument (following [16]) for the ‘‘almost orthogonality’’ of the DTTs is that they diagonalize the matrices in (30), which are almost symmetric and have pairwise distinct eigenvalues as mentioned above. For example, DCT-2 diagonalizes the symmetric $\phi(x)$ in (29) and hence can be made orthogonal by a suitable scaling $D \cdot DCT\text{-}2$, where D is diagonal.

F. Relationships Between DTTs

Some DTTs can be translated into each other using sparse matrices. These relationships can be understood and derived once their underlying signal models are known. We explained this in [9] (without using the notion of signal model) and briefly restate the result for completeness. The origin of these relationships is similarity in the signal model, i.e., that two DTTs belong to the same group of four (e.g., T -group).

Duality: We observed before that the right boundary conditions for the DTTs are precisely the mirrored versions of the left boundary conditions, a fact that meets our intuition since the DTTs are based on symmetric space models. However, the construction of $\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/p$ for a given DTT (see Theorem 6) deals differently with the left boundary condition (which determines the choice of the base sequence C) and the right boundary condition (which determines p); thus, we obtain different DTTs for a given pair of boundary conditions and for its mirrored counterpart. We call such a pair *dual*. Dual DTTs occur at mirrored positions in Table IV, i.e., at positions (i, j) ,

$(j, i), 1 \leq i, j \leq 4$, respectively. The DTTs on the main diagonal are self-dual.

Theorem 8 (Duality Relationship) [9]: Let DTT_n and DTT'_n be a pair of dual DTTs. Then

$$\text{diag}_{0 \leq k < n} ((-1)^k) \cdot \text{DTT}_n = \text{DTT}'_n \cdot J_n$$

where J_n is an identity matrix with the columns in reversed order. As an important consequence of Theorem 8, dual DTTs have the same arithmetic complexity.

Relationships in Groups of DTTs: Dual DTTs necessarily have the same associated $\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/p$. However, in Table IV, we also have DTTs that are not dual but have the same or similar \mathcal{M} , namely those in the same group of four (e.g., T -group). An example is given by the DCTs of type 3 and 4 with $\mathcal{M} = \mathbb{C}[x]/T_n$.

Further inspection shows that, in each group, all possible left and right boundary conditions are present. The DTTs in one group have (almost) the same module, but with different bases. Thus, we can translate DTTs in the same group into each other using a base change. Further, because of Table II, the resulting base change matrices are sparse, i.e., require only $O(n)$ operations.

Example: DCT, Type 3 and 4: We consider DCT-3_n and DCT-4_n , which are both in the T -group, i.e., the associated module is $\mathcal{M} = \mathbb{C}[x]/T_n$. The difference is in the choice of basis:

$$\begin{aligned} \text{DCT-3}_n &: \mathbb{C}[x]/T_n, \quad b = (T_0, \dots, T_{n-1}) \\ \text{DCT-4}_n &: \mathbb{C}[x]/T_n, \quad b' = (V_0, \dots, V_{n-1}). \end{aligned}$$

Using $T_\ell = (V_\ell + V_{\ell-1})/2$ from Table II and $V_{-1} = V_0$, the corresponding base change matrix S'_n for $b \rightarrow b'$ is given by

$$S'_n = \frac{1}{2} \cdot \begin{bmatrix} 2 & 1 & & & \\ 0 & 1 & 1 & & \\ & & \cdot & \cdot & \\ & & & 1 & 1 \\ & & & & 1 \end{bmatrix}. \quad (32)$$

We denote the zeros of T_n by $\alpha_k = \cos(k + 1/2)\pi/n$. As a consequence of the above, we get the following diagram:

$$\begin{array}{ccc} \mathbb{C}[x]/T_n & \xrightarrow{S'_n} & \mathbb{C}[x]/T_n \\ \downarrow \text{DCT-3}_n & & \downarrow \overline{\text{DCT-4}}_n \\ \bigoplus \mathbb{C}[x]/(x - \alpha_k) & \xrightarrow{I_n} & \bigoplus \mathbb{C}[x]/(x - \alpha_k) \end{array} \quad (33)$$

which implies the equation $\text{DCT-3}_n = \overline{\text{DCT-4}}_n \cdot S'_n$. Note that we have I_n in the bottom row of (33) since both $\text{DCT-3} = \overline{\text{DCT-3}}$ and $\overline{\text{DCT-4}}$ are polynomial transforms and thus use the same basis $(1, \dots, 1)$ in the spectrum. Introducing the scaling diagonal $D_n = \text{diag}_{0 \leq k < n} (\cos(2k + 1)\pi/(4n))$ of the DCT-4 (see Table IV), we get

$$D_n \cdot \text{DCT-3}_n = \text{DCT-4}_n \cdot S'_n. \quad (34)$$

If desired, this equation can now be further manipulated through transposition or inversion. As an example, one can obtain

$$S_n \cdot \text{DCT-2}_n \cdot \frac{1}{2} D_n^{-1} = \text{DCT-4}_n \quad (35)$$

where S_n is S'_n with the 2 replaced by 1 in the first entry and without the scaling factor 1/2.

Other Cases: Using this procedure on all DTTs shows that all DTTs of types 1–4 and all DTTs of type 5–8 can be converted into each other using $O(n)$ operations, respectively [9].

V. FINITE SKEW C -TRANSFORM AND SKEW DTTs

In this section, we introduce a new class of transforms that is closely related to the DTTs. We call these transforms *skew DTTs*. More specifically, the skew DTTs correspond to and generalize the DTTs in the T -group, i.e., those with associated $\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/T_n$, which are the DTTs of type 3 and 4. The first skew DCT (type 3) was introduced in [19].

We introduce the skew DTTs for the following reasons. First, they are interesting from a signal processing point of view. As the DTTs, they are associated with a finite space model, their associated boundary conditions are simple, and their signal extension is sparse even though not monomial.

Second, they are necessary building blocks in the general-radix Cooley–Tukey type DTT algorithms derived in [7].

A. Constructing the Signal Model

In the finite space models (24), we chose the right boundary condition to ensure a monomial signal extension via Lemma 5. Now we relax this requirement and consider a more general boundary condition for the four signal models in (24) for which $\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/T_n$. Namely we generalize to $\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/(T_n - \cos r\pi)$, $r \in \mathbb{Q}$, $0 \leq r \leq 1$. For $r = 1/2$, $\cos r\pi = 0$, which is the previous case. Hence, the Fourier transforms will generalize the DTTs in the T -group and depend on r .

The right boundary conditions associated with $\mathbb{C}[x]/(T_n - \cos r\pi)$ depend on the basis $C \in \{T, U, V, W\}$ and can be read off from Table IV:

$$\begin{aligned} T_n &= \cos r\pi \\ U_n &= U_{n-2} + 2 \cos r\pi \\ V_n &= -V_{n-1} + 2 \cos r\pi \\ W_n &= W_{n-1} - 2 \cos r\pi. \end{aligned} \quad (36)$$

In the general case $r \neq 1/2$, these boundary conditions lead to no monomial signal extensions, since this property uniquely defines the signal models for the 16 DTTs. However, it is intriguing that the signal extension is “two-monomial,” which means that the sum in (7) has at most two summands.

Lemma 9: The module $\mathbb{C}[x]/(T_n - \cos r\pi)$ with T -, U -, V -, or W -basis has a two-monomial signal extension.

Proof: The proof and the exact form of the signal extension can be found in [6]. ■

Resulting Finite Space Model: We define four skew finite space models parameterized by $r \in \mathbb{Q}$, $0 < r < 1$, for $V = \mathbb{C}^n$. Namely, for $C \in \{T, U, V, W\}$,

$$\begin{aligned} \mathcal{A} = \mathcal{M} &= \mathbb{C}[x]/(T_n - \cos r\pi) \\ \Phi : \mathbf{s} &\mapsto \sum_{0 \leq k < n} s_k C_k \in \mathcal{M}. \end{aligned} \quad (37)$$

As in (24), the natural basis of \mathcal{A} is the T -basis: $\mathcal{A} = \{h = \sum_{0 \leq k < n} h_k T_k\}$, independent of C . For $r = 1/2$, the skew models reduce to their nonskew counterpart in (24).

B. Spectrum and Fourier Transform: Skew DTTs

To compute the spectrum and a Fourier transform for the four models (37), we first need to determine the zeros of $T_n - \cos r\pi$ and fix a proper ordering.

Lemma 10: Let $r \in \mathbb{Q}$, $0 < r < 1$. We have the factorization

$$T_n - \cos r\pi = 2^{n-1} \prod_{0 \leq i < n} \left(x - \cos \frac{r+2i}{n} \pi \right) \quad (38)$$

which determines the zeros of $T_n - \cos r\pi$. We order the zeros as $\alpha = (\cos r_0\pi, \dots, \cos r_{n-1}\pi)$, such that $0 \leq r_i \leq 1$, and $r_i < r_j$ for $i < j$. The list α is given by the concatenation

$$\alpha = \bigcup_{0 \leq i < n/2} \left(\cos \frac{r+2i}{n} \pi, \cos \frac{2-r+2i}{n} \pi \right)$$

for even n , and by

$$\alpha = \left(\bigcup_{0 \leq i < \frac{n-1}{2}} \left(\cos \frac{r+2i}{n} \pi, \cos \frac{2-r+2i}{n} \pi \right) \right) \cup \left(\cos \frac{r+n-1}{n} \pi \right)$$

for odd n . In the particular case of $r = 1/2$ or $\cos r\pi = 0$, we thus have $\alpha = (\cos(i+1/2)\pi/n | 0 \leq i < n)$ as in Table VII.

Proof: The zeros of $T_n - \cos r\pi$ are proved using the closed form of T_n in Table VII. The ordering of α is shown by inspection. We omit the details. ■

In words, the list α arises from the list $\gamma = (\cos(r+2i)\pi/n | 0 \leq i < n)$ in (38) by interleaving the first half of γ with the reversed second half of γ .

Lemma 10 yields the Fourier transform for the models (37). We omit the form (8) and give directly the matrix forms \mathcal{F} .

Definition 11 (Skew DTTs): Let $p = T_n - \cos r\pi$, $0 < r < 1$, and $\mathcal{A} = \mathcal{M} = \mathbb{C}[x]/p$ with basis $b = (C_0, \dots, C_{n-1})$, where C is one of T, U, V, W . Let $\alpha = (\cos r_i\pi)_{0 \leq i < n}$ denote the list of zeros of p in the order specified in Lemma 10. We denote the associated polynomial transforms $\mathcal{P}_{b,\alpha}$ for \mathcal{M} by $\overline{\text{DCT-3}}_n(r)$, $\overline{\text{DST-3}}_n(r)$, $\overline{\text{DCT-4}}_n(r)$, $\overline{\text{DST-4}}_n(r)$, for $C = T, U, V, W$, respectively. Further, we define for each of these four $\overline{\text{DTT}}(r)$ the associated *scaled* polynomial transforms

$$\text{DTT}_n(r) = \text{diag}_{0 \leq i < n} (f(\cos r_i\pi)) \cdot \overline{\text{DTT}}_n(r)$$

where f is the scaling function associated with the (ordinary) DTT (see Table IV). We call these transforms *skew DTTs*. If $r = 1/2$, then $\overline{\text{DTT}}_n(1/2) = \overline{\text{DTT}}_n$ and $\text{DTT}_n(1/2) = \text{DTT}_n$ in all four cases. In the case of the $\text{DCT-3}_n(r) = \overline{\text{DCT-3}}_n(r)$, we will omit the bar for the skew versions. Specifically

$$\begin{aligned} \text{DCT-3}_n(r) &= [\cos r_k \ell \pi]_{0 \leq k, \ell < n} \\ \text{DST-3}_n(r) &= [\sin r_k (\ell + 1) \pi]_{0 \leq k, \ell < n} \\ \text{DCT-4}_n(r) &= [\cos r_k (\ell + 1/2) \pi]_{0 \leq k, \ell < n} \\ \text{DST-4}_n(r) &= [\sin r_k (\ell + 1/2) \pi]_{0 \leq k, \ell < n} \end{aligned}$$

As an example, we consider the $\text{DCT-4}_3(1/3)$. Using Lemma 10, the zeros of $T_3 - \cos(\pi/3) = T_3 - 1/2$ are given by $\alpha = (\cos(\pi/9), \cos(5\pi/9), \cos(7\pi/9))$. We get

$$\text{DCT-4}_3(1/3) = \begin{bmatrix} \cos \frac{1}{18} \pi & \cos \frac{1}{6} \pi & \cos \frac{5}{18} \pi \\ \cos \frac{5}{18} \pi & \cos \frac{5}{6} \pi & \cos \frac{11}{18} \pi \\ \cos \frac{7}{18} \pi & \cos \frac{5}{6} \pi & \cos \frac{1}{18} \pi \end{bmatrix}.$$

C. Filtering and Diagonalization Property

Filtering in the models (37) is multiplication of polynomials $h \in \mathcal{A}$, $s \in \mathcal{M}$ modulo $p = T_n - \cos r\pi$. In coordinates, it becomes the matrix-vector multiplication $\phi(h)\mathbf{s}$, where ϕ is the representation associated by the respective model. Convolution theorems are special cases of (11).

As an example, we compute the shift matrix $\phi(x)$ from (30) and (36). Specifically, it is obtained from (30) by adding in the upper right corner $\beta(r) = \cos r\pi$ for $\text{DCT-3}(r)$, and $\beta(r) = 2 \cos r\pi$ for the other skew transforms. Hence,

$$\phi(x) = \frac{1}{2} \cdot \begin{bmatrix} \beta_1 & 1 & & & \beta(r) \\ \beta_2 & 0 & 1 & & \\ 0 & 1 & 0 & \cdot & \\ & & 1 & \cdot & 1 \\ & & & \cdot & 0 & \beta_3 \\ & & & & 1 & \beta_4 \end{bmatrix}. \quad (39)$$

The values for the β_i coincide with the non-skew cases given in Table V. As a consequence, in the four cases

$$\text{DTT}(r) \cdot \phi(x) \cdot \text{DTT}(r)^{-1} = \text{diag}(\alpha)$$

where α is the list of zeros of $T_n - \cos r\pi$ from Lemma 10.

D. Translation Into Non-Skew DTTs

Each of the skew DTTs can be translated into its non-skew counterpart using a sparse x-shaped matrix.

Lemma 12: Let $\text{DTT}_n(r)$ be a skew DTT. Then

$$\text{DTT}_n(r) = \text{DTT}_n \cdot X_n^{(*)}(r)$$

and

$$\overline{\text{DTT}}_n(r) = \overline{\text{DTT}}_n \cdot X_n^{(*)}(r).$$

Here, $X_n^{(*)}(r)$ depends on the DTT and takes the following forms, indicated by $* \in \{C3, S3, C4, S4\}$. In all four cases, if the lines intersect, the numbers are added at the intersecting position.

$$\begin{aligned} X_n^{(C3)}(r) &= \begin{bmatrix} 1 & 0 & \cdots & \cdots & 0 \\ 0 & c_1 & & & s_{n-1} \\ \vdots & & \ddots & \ddots & \\ \vdots & & & \ddots & \\ 0 & s_1 & & & c_{n-1} \end{bmatrix} \\ X_n^{(S3)}(r) &= \begin{bmatrix} c_1 & & & & -s_{n-1} & 0 \\ & \ddots & \ddots & & & \vdots \\ & & \ddots & \ddots & & \vdots \\ -s_1 & & & & c_{n-1} & 0 \\ 0 & \cdots & \cdots & & 0 & c_n \end{bmatrix} \end{aligned}$$

with $c_\ell = \cos(1/2 - r)\ell\pi/n$ and $s_\ell = \sin(1/2 - r)\ell\pi/n$.

$$X_n^{(C4)}(r) = \begin{bmatrix} c'_0 & & & & s'_{n-1} \\ & \ddots & & & \\ & & \ddots & & \\ & & & \ddots & \\ s'_0 & & & & c'_{n-1} \end{bmatrix}$$

with $c'_\ell = \cos(1/2 - r)(2\ell + 1)\pi/(2n)$ and $s'_\ell = \sin(1/2 - r)(2\ell + 1)\pi/(2n)$. For DST-4(r), the sines s'_ℓ in $X_n^{(C4)}(r)$ are multiplied by -1 .

Proof: Follows by direct computation, using the definitions of the matrices and $\cos(x)\cos(y) = (\cos(x + y) + \cos(x - y))/2$. ■

Note that the 2×2 blocks in the translation matrices $X_n(r)$ are not rotations. The identities in Lemma 12 enable the inversion of the skew DTTs through the inversion of the ordinary DTTs.

E. Relationships Between Skew DTTs

All skew DTT(r) share the same associated module, but different bases. Thus, they can be translated into each other by a base change similar to the ordinary DTTs in Section IV-F. As in that section, we consider the skew DCTs, type 3 and 4, as an example. The base change matrix S'_n we computed in (32) did not depend on the right boundary condition. Thus, the diagram (33) generalizes for arbitrary r to

$$\begin{array}{ccc} \mathbb{C}[x]/(T_n - \cos r\pi) & \xrightarrow{S'_n} & \mathbb{C}[x]/(T_n - \cos r\pi) \\ \downarrow \text{DCT-3}_n(r) & & \downarrow \overline{\text{DCT-4}}_n(r) \\ \bigoplus \mathbb{C}[x]/(x - \alpha_k) & \xrightarrow{I_n} & \bigoplus \mathbb{C}[x]/(x - \alpha_k) \end{array} \quad (40)$$

which implies $\text{DCT-3}_n(r) = \overline{\text{DCT-4}}_n(r) \cdot S'_n$. The first difference occurs when we extend (40) to the non-polynomial $\text{DCT-4}_n(r)$, since the scaling diagonal depends on r . Let $\alpha = (\alpha_0, \dots, \alpha_{n-1})$ denote the zeros of $T_n - \cos r\pi$ and f the scaling function of DCT-4 and let $D_n(r) = \text{diag}_{0 \leq k < n}(f(\alpha_k))$. Then

$$D_n(r) \cdot \text{DCT-3}_n(r) = \text{DCT-4}_n(r) \cdot S'_n, \quad (41)$$

which generalizes (34).

In Section IV-F, we continued by inverting this equation to derive the different relationship (35). To do this, we introduce the proper “inverse” skew DTTs, which will also be needed in the DTT algorithms derived in [7]. The definition is motivated by and a generalization of the equations

$$\begin{aligned} \text{DCT-3}_n^{-1} &= 2/n \cdot \text{diag}(1/2, 1, \dots, 1) \cdot \text{DCT-2}_n \\ \text{DST-3}_n^{-1} &= 2/n \cdot \text{diag}(1, 1, \dots, 1/2) \cdot \text{DST-2}_n \\ \text{DTT}_n^{-1} &= n/2 \cdot \text{DTT}_n^T = n/2 \cdot \text{DTT}_n \end{aligned}$$

for DTT = DCT-4, DST-4.

Definition 13 (Inverse Skew DTTs): We define the *inverse skew DTTs* by

$$\begin{aligned} \text{idCT-3}_n(r) &= n/2 \cdot \text{diag}(2, 1, \dots, 1) \cdot \text{DCT-3}_n(r)^{-1} \\ \text{idCT-4}_n(r) &= n/2 \cdot \text{DCT-4}_n(r)^{-1} \\ \text{idST-3}_n(r) &= n/2 \cdot \text{diag}(1, 1, \dots, 2) \cdot \text{DST-3}_n(r)^{-1} \\ \text{idST-4}_n(r) &= n/2 \cdot \text{DST-4}_n(r)^{-1}. \end{aligned}$$

TABLE VI

OVERVIEW OF THE FINITE SPACE MODELS AND ASSOCIATED FOURIER TRANSFORMS DISCUSSED IN THIS PAPER

$\mathcal{A} = \mathcal{M}$	Φ	$\mathcal{F} = \mathcal{P}_{b,\alpha}$	other \mathcal{F}
$\mathbb{C}[x]/(x^2 - 1)U_{n-2}$	$\mathbf{s} \mapsto \sum s_k T_k$	DCT-1 $_n$	—
$\mathbb{C}[x]/T_n$		DCT-3 $_n$	—
$\mathbb{C}[x]/(x - 1)W_{n-1}$		DCT-5 $_n$	—
$\mathbb{C}[x]/(x + 1)V_{n-1}$		DCT-7 $_n$	—
$\mathbb{C}[x]/(T_n - \cos r\pi)$		DCT-3 $_n(r)$	—
$\mathbb{C}[x]/T_n$	$\mathbf{s} \mapsto \sum s_k U_k$	$\overline{\text{DST-3}}_n$	DST-3 $_n$
$\mathbb{C}[x]/U_n$		$\overline{\text{DST-1}}_n$	DST-1 $_n$
$\mathbb{C}[x]/V_n$		$\overline{\text{DCT-7}}_n$	DCT-7 $_n$
$\mathbb{C}[x]/W_n$		$\overline{\text{DST-5}}_n$	DST-5 $_n$
$\mathbb{C}[x]/(T_n - \cos r\pi)$		$\overline{\text{DST-3}}(r)_n$	DST-3 $_n(r)$
$\mathbb{C}[x]/(x - 1)W_{n-1}$	$\mathbf{s} \mapsto \sum s_k V_k$	$\overline{\text{DCT-6}}_n$	DCT-6 $_n$
$\mathbb{C}[x]/V_n$		$\overline{\text{DCT-8}}_n$	DCT-8 $_n$
$\mathbb{C}[x]/(x - 1)U_{n-1}$		$\overline{\text{DCT-2}}_n$	DCT-2 $_n$
$\mathbb{C}[x]/T_n$		$\overline{\text{DCT-4}}_n$	DCT-4 $_n$
$\mathbb{C}[x]/(T_n - \cos r\pi)$		$\overline{\text{DCT-4}}(r)_n$	DCT-4 $_n(r)$
$\mathbb{C}[x]/(x + 1)V_{n-1}$	$\mathbf{s} \mapsto \sum s_k W_k$	$\overline{\text{DST-8}}_n$	DST-8 $_n$
$\mathbb{C}[x]/W_n$		$\overline{\text{DST-6}}_n$	DST-6 $_n$
$\mathbb{C}[x]/T_n$		$\overline{\text{DCT-4}}_n$	DCT-4 $_n$
$\mathbb{C}[x]/(x + 1)U_{n-1}$		$\overline{\text{DST-2}}_n$	DST-2 $_n$
$\mathbb{C}[x]/(T_n - \cos r\pi)$		$\overline{\text{DST-4}}(r)_n$	DST-4 $_n(r)$

Thus, for $r = 1/2$, we have $\text{idCT-3}_n(1/2) = \text{DCT-2}_n$, $\text{idST-3}_n(1/2) = \text{DST-2}_n$, $\text{idCT-4}_n(1/2) = \text{DCT-4}_n$, $\text{idST-4}_n(1/2) = \text{DCT-4}_n$.

Note that Definition 13 does not provide direct knowledge about the matrix entries of the iDTTs. These, however, can be computed using Lemma 12. For example

$$\begin{aligned} \text{idCT-3}_n(r) &= \left(X_n^{(C3)}(r)\right)^{-1} \cdot \text{DCT-2}_n \\ \text{idCT-4}_n(r) &= \left(X_n^{(C4)}(r)\right)^{-1} \cdot \text{DCT-4}_n \end{aligned} \quad (42)$$

and similarly for DST-3 and DST-4. Note that $(X_n^{(*)}(r))^{-1}$ has in all four cases the same x-shaped pattern as $X_n(r)$. Namely, the four inverses are derived from

$$\begin{bmatrix} \cos a & \sin b \\ \sin a & \cos b \end{bmatrix}^{-1} = \frac{1}{\cos(a + b)} \begin{bmatrix} \cos b & -\sin b \\ -\sin a & \cos a \end{bmatrix}.$$

For example

$$\begin{aligned} \left(X_n^{(C3)}(r)\right)^{-1} &= \frac{1}{\cos(1/2 - r)\pi} \\ &\times \begin{bmatrix} c_n & 0 & \cdots & \cdots & 0 \\ 0 & c_{n-1} & & & -s_{n-1} \\ \vdots & & \ddots & \ddots & \\ \vdots & & & \ddots & \ddots \\ 0 & -s_1 & & & c_1 \end{bmatrix}. \end{aligned}$$

Using Definition 13, we can now invert (34) to get a generalization of (35)

$$S_n \cdot \text{idCT-3}_n(r) \cdot \frac{1}{2} D_n(r)^{-1} = \text{idCT-4}_n(r)_n$$

where S_n is the same as in (35).

TABLE VII
FOUR SERIES OF CHEBYSHEV POLYNOMIALS. IN THE TRIGONOMETRIC CLOSED FORM $\cos \theta = x$ AND IN THE POWER FORM $(u + u^{-1})/2 = x$

C_n	$n = 0, 1$	closed form	power form	symmetry	zeros ($0 \leq k < n$)
T_n	$1, x$	$\cos(n\theta)$	$\frac{u^n + u^{-n}}{2}$	$T_{-n} = T_n$	$\cos \frac{(k+\frac{1}{2})\pi}{n}$
U_n	$1, 2x$	$\frac{\sin(n+1)\theta}{\sin \theta}$	$\frac{u^{n+1} - u^{-(n+1)}}{u - u^{-1}}$	$U_{-n} = -U_{n-2}$	$\cos \frac{(k+1)\pi}{n+1}$
V_n	$1, 2x - 1$	$\frac{\cos(n+\frac{1}{2})\theta}{\cos \frac{1}{2}\theta}$	$\frac{u^{n+1/2} + u^{-(n+1/2)}}{u^{1/2} + u^{-1/2}}$	$V_{-n} = V_{n-1}$	$\cos \frac{(k+\frac{1}{2})\pi}{n+\frac{1}{2}}$
W_n	$1, 2x + 1$	$\frac{\sin(n+\frac{1}{2})\theta}{\sin \frac{1}{2}\theta}$	$\frac{u^{n+1/2} - u^{-(n+1/2)}}{u^{1/2} - u^{-1/2}}$	$W_{-n} = -W_{n-1}$	$\cos \frac{(k+1)\pi}{n+\frac{1}{2}}$

VI. OVERVIEW OF FINITE SPACE MODELS

In Table VI, we list all the finite space signal models, and their associated Fourier transforms, that we introduced in this paper. The table is divided according to Φ , which is a finite T -, U -, V -, or W -transform.

In each row, we list in the first two columns the signal model, in the third column the associated unique polynomial Fourier transform, and in the fourth column possibly other relevant Fourier transforms for the model.

Except for the skew DTTs, each of the listed transforms has an orthogonal counterpart, which is obtained by proper scaling of rows or columns.

Table VI, together with Table II in [1] for finite 1-D time models classifies practically all existing 1-D trigonometric transforms, i.e., those transforms that can be expressed using cosines and sines. For each of these transforms, ASP hence provides the associated signal model and with it all basic SP concepts, many of which have not been defined or found before.

VII. CONCLUSION

This paper shows that a theory of linear signal processing can be developed from a new concept of shift that is different from the standard time shift, namely from the space shift as we call it. Using the algebraic signal processing theory, we derived from this shift appropriate signal models for space signal processing, i.e., filter algebras, signal modules, and “ z -transforms.” In the finite case this approach derived from basic principles the 16 DTTs as Fourier transforms. This interpretation is arguably more satisfying than the original one as asymptotic approximations of the Karhunen–Loève transform (KLT) of a first-order causal Gauss–Markov random process. For a closer investigation of the relationship between KLTs and DTTs and between KLTs and general Fourier transforms in ASP see [6], [20].

By identifying the signal models underlying the DTTs, we also identified their associated notions of “ z -transform,” filtering or convolution, and explained in one framework many of the known properties of the DTTs. In [7], [9] we use the knowledge of these signal models to derive known and novel fast DTT algorithms.

One may wonder which other shifts provide meaningful SP frameworks and ASP is the proper platform to investigate this question. We have done first steps in this direction with a generalization of the space shift (called GNN shift) in [6], and with 2-D space shifts for both the quincunx lattice [21] and the hexagonal lattice [22]. The latter two yield nonseparable 2-D signal models.

APPENDIX I

CHEBYSHEV POLYNOMIALS

Chebyshev polynomials, and the more general class of orthogonal polynomials, have many interesting properties and play an important role in different areas of mathematics, including statistics, approximation theory, and graph theory. An excellent introduction to the theory of orthogonal polynomials can be found in the books of Chihara, Szegő, and Rivlin [23]–[25]. In this section we give the main properties of Chebyshev polynomials that we will use in this paper.

We call every sequence $C = (C_n)_{n \in \mathbb{Z}}$ of polynomials that satisfies the three-term recurrence

$$C_{n+1}(x) = 2xC_n(x) - C_{n-1}(x) \tag{43}$$

a sequence of *Chebyshev polynomials* (C stands for Chebyshev). Using (43), the sequence C is uniquely determined by the initial polynomials C_0, C_1 . The most important—and commonly known—are the Chebyshev polynomials of the *first kind*, denoted by $C_n = T_n$ and determined by $T_0 = 1$ and $T_1 = x$. We provide a few examples:

$$\frac{T_{-2} \quad T_{-1} \quad T_0 \quad T_1 \quad T_2 \quad T_3}{2x^2 - 1 \quad x \quad 1 \quad x \quad 2x^2 - 1 \quad 4x^3 - 3x}$$

For $x \in [-1, 1]$, T_n can be written in closed form as

$$T_n = \cos n\theta, \quad \cos \theta = x. \tag{44}$$

The closed form exhibits the *symmetry property* $T_{-n} = T_n$, $n \in \mathbb{Z}$, and can be used to derive the zeros of T_n . We will occasionally use another parameterization of T_n , which we call *power form*, given by

$$T_n = \frac{u^n + u^{-n}}{2}, \quad \frac{u + u^{-1}}{2} = x. \tag{45}$$

By substituting $u = e^{j\theta}$ we obtain (44).

In this paper, we also consider the Chebyshev polynomials of the second, third, and fourth kind, denoted by U_n, V_n, W_n , respectively, that arise from $C_0 = 1$ and different choices of C_1 . Each of these sequences exhibits a symmetry property and possesses parameterized forms. These properties are summarized in Table VII.

In addition, we will need the following properties that are shared by all sequences of Chebyshev polynomials including T, U, V, W (see [23]).

Lemma 14: Let $C = (C_n)_{n \in \mathbb{Z}}$ be a sequence of Chebyshev polynomials. Then the following holds:

- i) the sequence C is determined by any two successive polynomials C_n, C_{n+1} ;
- ii) $\deg(C_0) = 0, \deg(C_1) = 1 \Rightarrow \deg(C_n) = n, \text{ for } n \geq 0$;
- iii) $C_n = C_1 \cdot U_{n-1} - C_0 \cdot U_{n-2}$;
- iv) $T_k \cdot C_n = (C_{n+k} + C_{n-k})/2$.

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